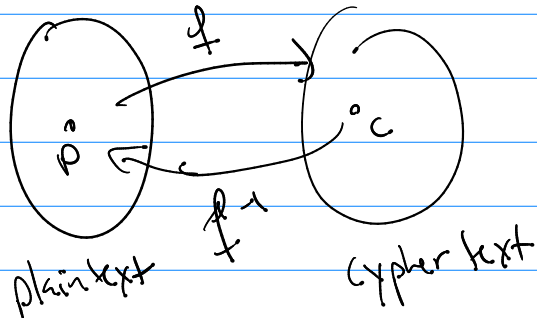


Math 321

Cypho



f : encryption func
 f^{-1} : decryption fun

Private Key

→ Character Cypher

plaintext: $p_1 p_2 p_3 p_4 \dots p_n$ are characters

→ $n = | \text{characters of language} |$

cypher text is usually the same characters

⊛ f is just a bijection from char. to char.

⊙ shift $f(p) = (p + k) \text{ mod } n$

Visual

as
 $\begin{matrix} a & b \\ p & c \\ e & a \end{matrix}$

$$f(p) = (p + 2) \text{ mod } 7$$

$$dad \rightarrow fc f$$

$$f^{-1}(c) = (c - 2) \text{ mod } 7$$

+5

Note

Cryptanalysis : by study of cypher text

(with "knowing" f), Find plaintext?

(2) affine shift

$$f(p) = (ap + k) \pmod n$$

$$f^{-1}(c) = (\bar{a}(c - k)) \pmod n$$

$\rightarrow \bar{a}$ is a 's inv. mod n

So $\gcd(a, n) = 1$ is required!

ex charc: {a, b, c, d, e, f, g} (charc = 7)

$$f(p) = (5p + 3) \pmod 7$$

ex 'd' $\rightarrow f(3) = (5 \cdot 3 + 3) \pmod 7 = 4 \rightarrow$ 'e'
'a' $\rightarrow f(0) = (5 \cdot 0 + 3) \pmod 7 = 3 \rightarrow$ 'd'

dad \rightarrow ede

and $f^{-1}(c) = (\bar{5}(c - 3)) \pmod 7$

Find $\bar{5} \pmod 7$: $\gcd(7, 5)$ $7 = (1)5 + (2) \leftarrow p_2$
 $\gcd(5, 2)$ $5 = (2)2 + (1) \leftarrow p_1$
 $\gcd(2, 1)$ $2 = (2)1 + 0$

by P_1 $\boxed{1} = 5 + (-2) \boxed{2}$
 by P_2 $1 = 5 + (-2) \boxed{7 + (-1)5}$

$$1 = \underset{\uparrow}{(3)} 5 + \underset{\downarrow}{(-2)} 7$$

$$\boxed{3} = \overline{5} \pmod{7} \quad -2 = \overline{7} \pmod{5}$$

PK in our example

$$f^{-1}(c) = (3(c-3)) \pmod{7}$$

(3) random 1-1 replacement

$$a \leftrightarrow b$$

$$b \leftrightarrow g$$

$$c \leftrightarrow c$$

$$d \leftrightarrow e$$

$$e \leftrightarrow a$$

$$f \leftrightarrow d$$

$$g \leftrightarrow f$$

$$dad \rightarrow ebe$$

(4) one-time-pad

Plain text $P_1 P_2 P_3 \dots P_n$

Random #'s $r_1 r_2 r_3 \dots r_n$

↓

Cypher: $C_i = (P_i + r_i) \pmod{n}$

$C_1 C_2 C_3 \dots C_n$

$$f^{-1}(c_i) = (c_i - r_i) \bmod n$$

Public Key

RSA, Cocks cryptography

① p, q are large primes (these are the true private key)

② $n = pq$ $M = (p-1)(q-1)$

③ pick a number, e , so that

④ find $d = e^{-1} \bmod M$ } exists b/c rel. prime

Public (e, n) private (d, p, q, k)
public key

$$f(p) = p^e \bmod n$$

$$f^{-1}(c) = c^d \bmod n$$

(RSA)

Do Cocks Crypto

$$p = 7 \quad q = 13 \quad n = 91 \quad m = 72$$

$$\rightarrow e \text{ so } \gcd(e, n) = 1 \quad \text{let } e = 5$$

$d = \bar{e}$
2

$$\gcd(72, 5) \quad 72 = (14)5 + \textcircled{2}$$

$$5 = (2)2 + \boxed{1}$$

$$\boxed{1} = (1)5 + (-2)\textcircled{2}$$

$$1 = (1)5 + (-2)(72 + (-14)5)$$

$$1 = (29)5 + (-2)72$$

$$\bar{5} = \underline{29} \pmod{72}$$

$$\begin{aligned} f(p) &= p^5 \pmod{91} \\ f^{-1}(c) &= c^{29} \pmod{91} \end{aligned}$$

to use this Note: $b(c \pmod{91})$ we can have
at most 91 symbols.

if \pmod{n} is big ..

(ex)

$\pmod{19,519}$

Mark

1
1200

alt(w)

$$\gcd(149, k)$$

$$\rightarrow \underline{\underline{12 \pmod{149}}} \quad (\text{as pos. int})$$