

Math 321

Q's /

5.2 #11

Nim

1, 2, 3, ..., n

n - matches.

each time you play to take 1, 2, or 3. Take all

that is left \rightarrow you lose.

player 1 sees

$$n = 4j \text{ to}$$

$$n = 4j + 2$$

$$n = 4j + 3$$

Player 1 wins

player 1 sees

$$n = 4j + 1$$

Player 1 loses

Player 1 loses

if $n \bmod 4 \rightarrow 0, 2, 3 \rightarrow$ win
 $n \bmod 4 = 1 \rightarrow$ lose

play:

$n = 1 \rightarrow$ player 1 takes it and loses

$n = 2 \rightarrow$ player 1 takes 1 \rightarrow player 2 faces $n = 1$

\therefore player 1 wins

$n = 3 \rightarrow$ player 1 takes 2 \rightarrow player 2 gets $n = 1$

\therefore player 1 wins

$n = 4 \rightarrow$ player 1 takes 3 \rightarrow player 2 gets $n = 1$

\therefore player 1 wins

$n = 5 \rightarrow$ 3 cases

player 1 takes 1 \rightarrow p2 gets 4 $\rightarrow \therefore$ p1 lose

player 1 takes 2 \rightarrow p2 gets 3 \rightarrow p1 lose

player 1 takes 3 \rightarrow p2 gets 2 \rightarrow p1 lose

Show: Win by take 1, 2, or 3

$$n \bmod 4 = 1 \rightarrow \text{player 1 loses}$$
$$0, 2, 3 \rightarrow \text{player 1 wins.}$$

Def

Base:

show formula works for $n=1, 2, 3, 4$

Inductive (strong)

Assume $P(n=1) \wedge P(n=2) \wedge \dots \wedge P(n=k)$

show $P(n=k+1)$

$$n = k+1 \rightarrow (k+1) \bmod 4$$

\rightarrow Case #1 $(k+1) \bmod 4 = \underline{0}$

so player 1 takes 3 $(k-2) \bmod 4 = 1$

$$\rightarrow \text{player 2 sees } \underline{\underline{(k+1) - 3}} \bmod 4$$

$$= (-3) \bmod 4 = 1$$

$\therefore p_2$ loses

Case #2 $(k+1) \bmod 4 = \underline{1}$

(2a) take 1 $\underline{\underline{(k+1) - 1}} \bmod 4 = \underline{0}$

(2b) take 2 $\underline{\underline{(k+1) - 2}} \bmod 4 = \underline{3}$

(2c) take 3 $\underline{\underline{(k+1) - 3}} \bmod 4 = \underline{2}$

p_1 wins
 $\rightarrow p_1$ loses

5.1 #2a

H_n is the n^{th} harmonic number

$$H_k = 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{k}$$

$$H_{2^n} \leq 1 + n$$

$$1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \dots + \frac{1}{2^n} \leq 1 + n$$

$$n=0 \quad H_{2^0} = H_1 = 1$$

$$n=1 \quad H_2 = 1 + \frac{1}{2}$$

$$n=2 \quad H_4 = 1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4}$$

$$n=3 \quad H_8 = 1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \frac{1}{5} + \frac{1}{6} + \frac{1}{7} + \frac{1}{8}$$

⋮
show $H_{2^n} \leq 1 + n$

DB **Base** $P(1^{\text{st}} \text{ case}) = P(n=0) : "H_1 \leq 1 + 0"$

$$= "1 \leq 1" \quad \text{true}$$

Inductive

Assume $P(k^{\text{th}})$
 $n=k$

$$H_k \leq 1 + k$$

$1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{2^k}$

I.H.

Show: $P(k+1^{\text{st}})$: $H_{k+1} \leq 1 + (k+1)$
 $n=k+1$

$$1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{2^{k+1}} \leq \dots$$

$$\text{task } H_{2^k} = \overbrace{\left(1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{2^k}\right)}^{2^k \text{ terms}} \overbrace{\left(\frac{1}{2^{k+1}} + \dots + \frac{1}{2^{k+1}}\right)}^{2^k \text{ terms}}$$

$$\leq \underbrace{(1+k)}_{1^k} + \underbrace{\frac{1}{2^{k+1}}}_{\frac{1}{2^k}} + \dots + \underbrace{\frac{1}{2^{k+1}}}_{\frac{1}{2^k}} \leq \frac{1}{2^k}$$

$$\leq (1+k) + \frac{1}{2^k} + \frac{1}{2^k} + \dots + \frac{1}{2^k}$$

$$\leq (1+k) + 2^k \left(\frac{1}{2^k}\right) = 1+k+1 \checkmark$$

Exan 3 11 probs

4.1 $a|b, a \text{ mod } m, a \equiv b \pmod{m}$ (2 probs)

(1) divisibility proof

ex $a|b \wedge b|c \rightarrow a|c$

$ak = b$ for an integer k

(2) find $div, mod, a = qb + r$

ex $-21 \text{ mod } 4, 21 \text{ div } 6,$
 $(32^{121})^{1021} \text{ mod } 3)$

4.2 Int. Reps and ops (1 prob)

(1) $()_b$ $()_b$ $b = 2, 5 \text{ or } 7$
 $+ ()_b$ $\times ()_b$

4.3 Primes, GCD, LCM, Euclid's (3 probs)

- ① Prove Primes are infinite
- ② gcd, lcm using prime factors.
- ③ Euclid's Alg. \Leftrightarrow Bézout's th²
gcd(a,b) by euclid's

$$\text{show } \text{gcd}(a,b) = sa + tb$$

4.6 Crypto (2 probs)

- ① Shift, Affine shift, ^{and} or One-time pad

plaintext $\xrightarrow{\quad}$ cypher text

- ② for RSA/locks crypto.

given $[e, n]$ \rightarrow write the encryption
and decryption functions.

$$c = p^e \pmod n$$

$$p = c^d \pmod n$$

5.1-5.3

3 probs

Induction

(1) Weak Ind, proof

(2) Strong Ind, proof

(3) Weak Ind, proof involving f_0, f_1, \dots