

Math 321

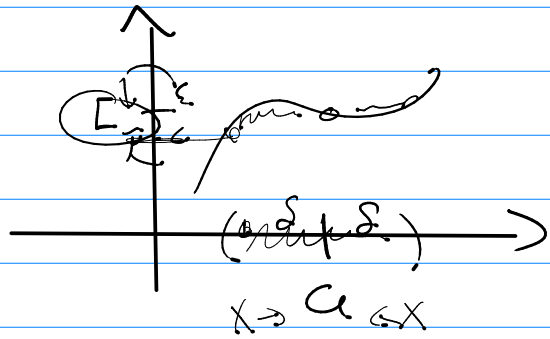
Quantification

Proportional function

$$\lim_{x \rightarrow a} f(x) = L$$

$$\forall \epsilon > 0 \exists \delta > 0$$

$$\text{if } 0 < |x - a| < \delta, \text{ then } |f(x) - L| < \epsilon$$



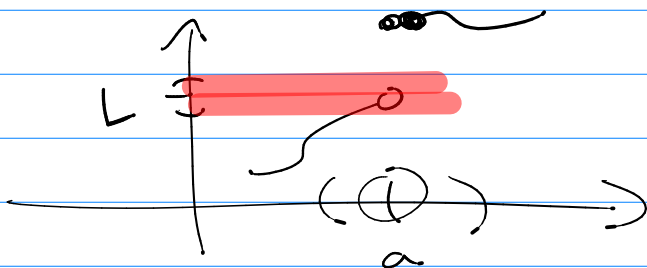
when do we say $\lim_{x \rightarrow a} f(x)$ d.u.e.?

(this is not true)

$$\neg (\forall \epsilon > 0 \exists \delta > 0 \ 0 < |x - a| < \delta \rightarrow |f(x) - L| < \epsilon)$$

$$\equiv \exists \epsilon > 0 \forall \delta > 0 \neg (0 < |x - a| < \delta \wedge |f(x) - L| < \epsilon)$$

$$\equiv \exists \epsilon > 0 \forall \delta > 0 (0 < |x - a| < \delta \wedge \neg (|f(x) - L| < \epsilon))$$



Argument

Seq of statements end with a conclusion.

↳ Premise.

(Statement \wedge statement \wedge .. \wedge statement) \rightarrow Conclusion

Valid if it is always true

D	Δ	$D \rightarrow \Delta$
T	T	T
T	F	F
F	T	T
F	F	T

If you have true statements force conclusion to be true \rightarrow tautology!

(Valid)

Notation:
Premise
Premise
:
Premise
 \therefore Conclusion

Argument Form.

Premises, conclusion have prop. variables

(ex) $[(p \rightarrow q) \wedge p] \rightarrow q$

$$\begin{array}{l} p \rightarrow q \\ p \\ \hline \therefore q \end{array}$$

if a tautology, it is Valid

$$\begin{aligned}
 [(p \rightarrow q) \wedge p] \rightarrow q &\equiv \neg [(p \rightarrow q) \wedge p] \vee q \\
 &\equiv [\neg(p \rightarrow q) \vee \neg p] \vee q \\
 &\equiv \neg(p \rightarrow q) \vee (\neg p \vee q) \equiv \neg(p \rightarrow q) \vee (p \rightarrow q) \\
 &\equiv T
 \end{aligned}$$

$$\begin{array}{l}
 p \rightarrow q \\
 p \\
 \hline
 \therefore q
 \end{array}$$

affirming the hypothesis / modus ponens

$$\begin{array}{l}
 p \rightarrow q \\
 \neg q \\
 \hline
 \therefore \neg p
 \end{array}$$

denying the conclusion / modus tollens

$$\begin{array}{l}
 p \rightarrow q \\
 q \rightarrow r \\
 \hline
 \therefore p \rightarrow r
 \end{array}$$

hypothetical
syllogism

$$\begin{array}{l}
 p \vee q \\
 \neg p \\
 \hline
 \therefore q
 \end{array}$$

disjunctive
syllogism

$$\begin{array}{l}
 p \\
 \hline
 \therefore p \vee q
 \end{array}$$

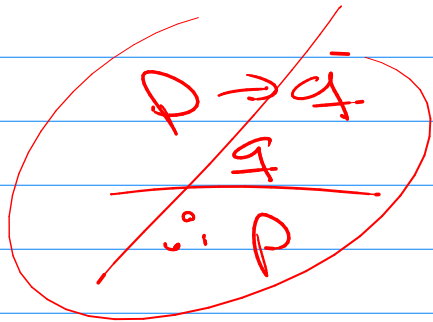
addition

$$\begin{array}{l}
 p \wedge q \\
 \hline
 \therefore p
 \end{array}$$

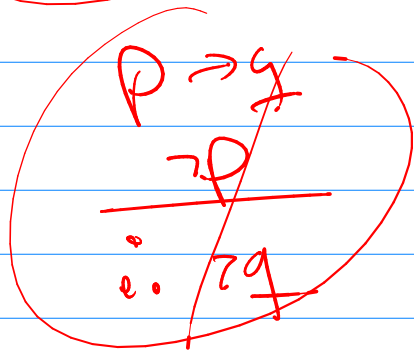
simplification

Note: use an Invalid argument form as
 if it was a rule of inference

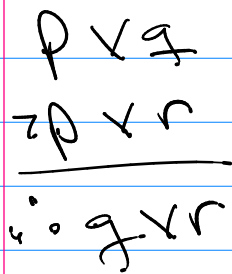
→ It is called a fallacy



fallacy of affirming the conclusion



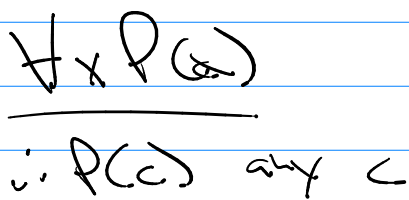
fallacy of denying the hypothesis



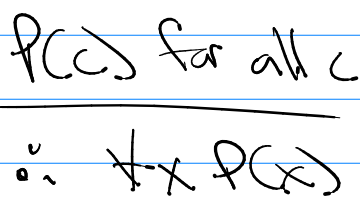
resolution

Quantification:

actual measures → general statements
 of all, some



univ. instantiation



univ. generalization

$\exists x P(x)$

$\therefore P(c)$ for some c

Existential Instantiation

$P(c)$ for some c

$\therefore \exists x P(x)$

Existential Generalization
