

Math 321

Q5 $p = 53$ $q = 61$ $e = 17$ $n = 3233$ $e = 17$
 $M = 3120$

d is the inverse mod $M = 3120$

$\gcd(3120, 17) =$

$3120 = 183 \cdot 17 + 9$

ex 8

A = 01 H = 08 O = 15 V = 22

B = 02 I = 09 P = 16 W = 23

C = 03 J = 10 Q = 17 X = 24

D = 04 K = 11 R = 18 Y = 25

E = 05 L = 12 S = 19 Z = 26

F = 06 M = 13 T = 20

G = 07 N = 14 U = 21

UPLOAD

start @ 1

UP = 2116

LO = 1245

AD = 0104 = 104

start @ 0

UP = 2015

LO = 1114

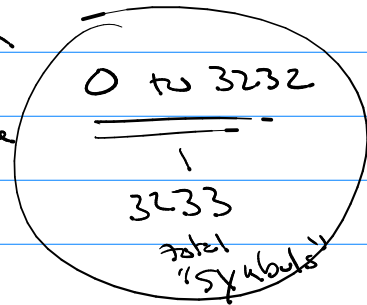
AD = 3

Block size:

AAA
:
zzz

$26^3 > 3233$ $26^2 < 3233$

$E(p) = p^{17} \pmod{3233}$ → return a value



$P_1 = 2015$

$C_1 = 2015^{17} \pmod{3233}$

decrypt $P_1 = C_1^d \pmod{3233}$

Ch 5 Induction

$n=2,3,4, \dots$

Q (if $n^2 < 2^n \rightarrow n > 4$) \equiv ($2 \leq n \leq 4 \rightarrow n^2 \geq 2^n$)
 $\left. \begin{matrix} n=2 \\ n=3 \\ n=4 \end{matrix} \right\}$ cases

conjecture:

$n=5, 6, 7, \dots \rightarrow n^2 < 2^n$

infinite cases

$\forall n P(n)$

where u.d. $n=5, 6, 7, \dots$
 $P(n)$: " $n^2 < 2^n$ "

Prop: Finite set for n $\forall n P(n) \equiv P(e_1) \wedge P(e_2) \wedge \dots \wedge P(e_k)$

infinite? $\forall n P(n) \equiv P(e_1) \wedge P(e_2) \wedge \dots$

Ch 5 Induction:

technique to prove $\forall n P(n)$ where u.d. is infinite and is a well-ordered set

$e_1 \leq e_2 \leq e_3 \leq e_4 \leq \dots$

basically there is always a 1st element for any subset of a well-ordered set.

Goal: $\forall n P(n)$ is true

- ① we will use two tactics to help prove this
- ② tactics are based on well-ordered sets
 set = $e_1 \leq e_2 \leq e_3 \leq e_4 \leq \dots$

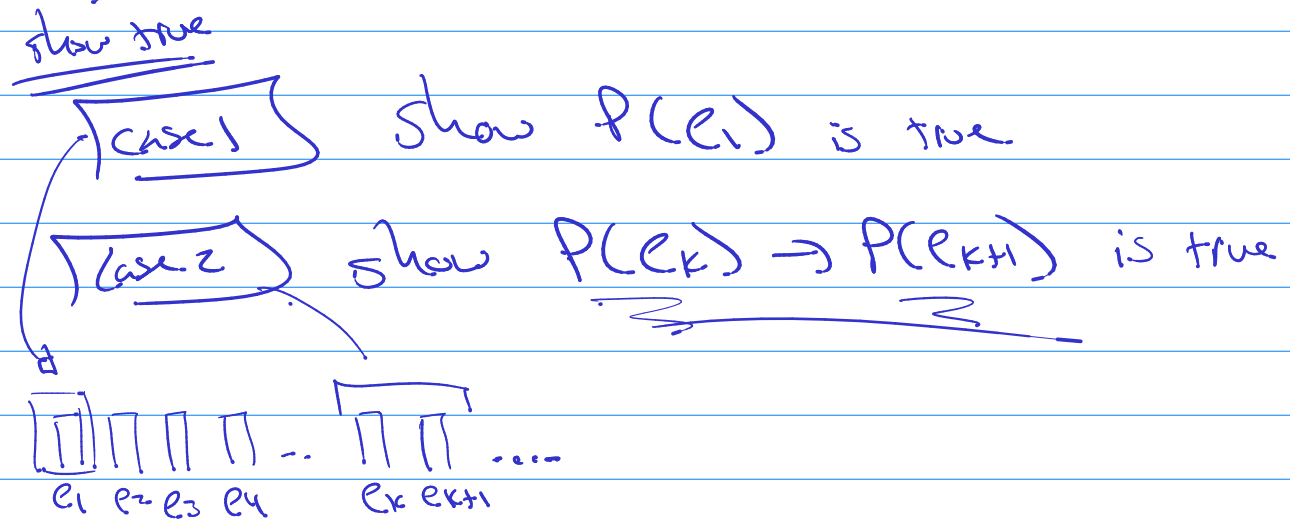
Tablelogy #1

$$\boxed{[P(e_1) \wedge \forall k (P(e_k) \rightarrow P(e_{k+1}))]} \rightarrow \underbrace{\forall n P(n)}_P \equiv \text{True}$$

L	R	L → R
T	T	T
T	F	F
F	T	F
F	F	F

→ $\forall n P(n)$ being true is nec. when left is true.

So prove the left $[P(e_1) \wedge \forall k (P(e_k) \rightarrow P(e_{k+1}))]$



Prove by "weak" induction $\forall n P(n)$, $n = e_1, e_2, e_3, \dots$

Case 1 (Basis Step) show $P(e_1)$ is true

Case 2 ("weak" Inductive Step) show $P(e_k) \rightarrow P(e_{k+1})$

2nd technology:

$[P(e_1) \wedge \forall k (P(e_1) \wedge P(e_2) \wedge \dots \wedge P(e_k) \rightarrow P(e_{k+1}))] \rightarrow \forall n P(n)$

Case 1 (Basis Step) show $P(e_1)$ is true

Case 2 (strong inductive step)

show $(P(e_1) \wedge P(e_2) \wedge \dots \wedge P(e_k)) \rightarrow P(e_{k+1})$ is true

Prf.: $\forall n$ " $1+2+\dots+n = \frac{n(n+1)}{2}$ ", $n=1, 2, 3, \dots$

$P(n)$: " $1+2+\dots+n = \frac{n(n+1)}{2}$ " $n=1, 2, 3, \dots$

Proof by Induction:

(Basis Step) show $P(e_1)$ is true

e_1 : $n=1$ $P(e_1)$: " $1 = \frac{1(1+1)}{2}$ " True!

(Inductive step) show $P(e_k) \rightarrow P(e_{k+1})$

$P(n=k)$: " $1+2+\dots+k = \frac{k(k+1)}{2}$ " $\overset{n=k}{\uparrow}$ $\overset{n=k+1}{\downarrow}$

$P(n=k+1)$: " $1+2+\dots+(k+1) = \frac{(k+1)(k+2)}{2}$ "