

Math 321

Ch 8 Back to inductive / recursive relations

8.13 why?

ex) $t_0 = 0, t_1 = 1$ ← Basis
 $f_n = f_{n-1} + f_{n-2} \quad n = 2, 3, 4, \dots$

A recursive / inductive relation

(relationship of old to new) → old elements come together to make new elements

→ this is a natural way to model.

ex) tower of hanoi



$$H_1 = 1$$

$$H_2 = H_1 + 1 + H_1 = 2H_1 + 1$$

$$H_3 = H_2 + 1 + H_2 = 2H_2 + 1$$

$$H_n = 2H_{n-1} + 1$$

Model

Basis: $H_1 = 1$

Recursive relation: $H_n = 2H_{n-1} + 1$

①

Seq: 1, 3, 7, 15, 31, 63, ...

$H_1, H_2, H_3, H_4, H_5, H_6, \dots$

② recursive \rightarrow closed formula function

Solving the recurrence relation

How? ex $H_n = 2H_{n-1} + 1, H_1 = 1$

tech #1 guess seq: 1, 3, 7, 15, 31, 63, ...

$$a_n = 2^n - 1 \quad n=1, 2, 3, \dots$$

check:

$$H_1 = 1 \quad a_1 = 2^1 - 1 = 1$$
$$H_n = 2H_{n-1} + 1 \rightarrow 2^n - 1 \stackrel{?}{=} 2(2^{n-1} - 1) + 1$$

$$2^n - 1 \stackrel{?}{=} 2^n - 2 + 1$$

true

tech #2 iteration (forward/backward)

$$H_1 = 1$$

$$H_n = 2 \cdot H_{n-1} + 1$$

$$H_1 = 1$$

$$H_2 = 2 \cdot 1 + 1 = 2 + 1$$

$$H_3 = 2(2+1) + 1 = 2^2 + 2 + 1$$

$$H_4 = 2(2^2 + 2 + 1) + 1 = 2^3 + 2^2 + 2 + 1$$

$$H_n = 2^{n-1} + 2^{n-2} + \dots + 2^1 + 2^0$$

$$H_n = \sum_{k=0}^{n-1} 2^k = \frac{2^n - 1}{2 - 1} = 2^n - 1$$

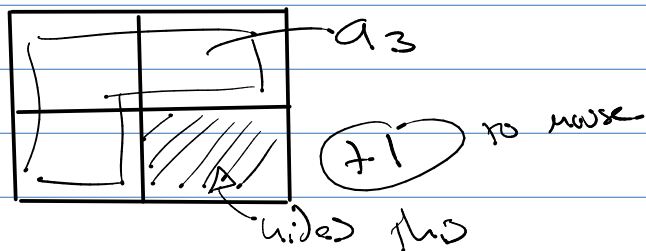
$$\sum_{k=1}^n a_k = a \left(\frac{r^{n+1} - 1}{r - 1} \right)$$

Modeling with rec. relations

ex 4 kids. Went around town to get all the cookies.
→ take them all home and share them. Wait one day

over night. Kid #4

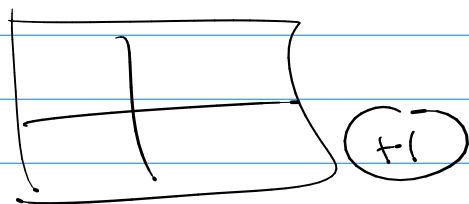
$$a_4 = \frac{4}{3} a_3 + 1$$



Kid #3 $a_3 = \frac{4}{3} a_2 + 1$

Kid #2 $a_2 = \frac{4}{3} a_1 + 1$

Kid #1 $a_1 = \frac{4}{3} a_0 + 1$



a_0 is what is there in the morning.





Notice: $4 \mid a_0$


Seq: a_0, a_1, a_2, a_3, a_4 ?


Basis: $4 \mid a_0$

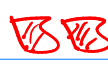
Inductive $a_n = \frac{4}{3} a_{n-1} + 1 \quad n=1, 2, 3, 4$

(ex)

how many tiles of length n used that are ,  never have a  in it?

[ex]  is ok (count)

 is not ok (do not count)

$C_n = \#$ of ways to tile n length w/o 

Base:

$$C_1 = 2$$

$$C_2 = 3$$









Inductive:

$$C_n = (\text{1 blue tile})_n (\text{rest}) \text{ or } (\text{red/blue})_n (\text{rest})$$

$$C_n = 1 \cdot C_{n-1} + 1 \cdot C_{n-2}$$

$$C_n = C_{n-1} + C_{n-2}$$

Mod:

$$\begin{cases} C_1 = 2 \\ C_2 = 3 \end{cases}$$

base

recursive

$$C_n = C_{n-1} + C_{n-2}$$

$$C_n = f(n)$$

seq:

2, 3, 5, 8, 13, ...