

Math 321

Q's

1.1 (15e)

ex \square is ^{and} nec ~~but~~ ~~(not)~~ suff for Δ
 $(\rightarrow) \quad \neg(\rightarrow)$

$\rightarrow \square$ is nec. for $\Delta \quad \Delta \rightarrow \square$ \star

$\rightarrow \square$ is suff. for $\Delta \quad \square \rightarrow \Delta$

$\rightarrow \square$ is not suff for Δ

it is not the case that \square is suff for $\Delta \quad \neg(\square \rightarrow \Delta)$
 \star

$$(\Delta \rightarrow \square) \wedge \neg(\Delta \rightarrow \Delta)$$

15e for hiking, it is nec but not suff that (\neg ripe \wedge \neg bears)

$$\left[\text{hiking} \rightarrow (\neg \text{ripe} \wedge \neg \text{bears}) \right] \wedge \neg \left[(\neg \text{ripe} \wedge \neg \text{bears}) \rightarrow \text{hiking} \right]$$

hiking	ripe	bears	$\neg \text{ripe} \wedge \neg \text{bears}$	$\text{hike} \rightarrow (\neg \text{ripe} \wedge \neg \text{bears})$
T	T	T	F	F
T	T	F	F	F
T	F	T	F	F
T	F	F	T	T
F	T	T	F	T
F	T	F	F	T
F	F	T	T	T
F	F	F	T	T

Distributive: $\neg(p \wedge q) \equiv \neg p \vee \neg q$ $\neg(p \vee q) \equiv \neg p \wedge \neg q$

$$\begin{aligned}
 p \rightarrow (q \wedge r) &\equiv (p \rightarrow q) \wedge (p \rightarrow r) \\
 p \rightarrow (q \vee r) &\equiv (p \rightarrow q) \vee (p \rightarrow r) \\
 p \wedge (q \vee r) &\equiv (p \wedge q) \vee (p \wedge r) \\
 p \vee (q \wedge r) &\equiv (p \vee q) \wedge (p \vee r) \\
 (q \wedge r) \rightarrow p &\equiv (q \rightarrow p) \wedge (r \rightarrow p) \\
 (q \vee r) \rightarrow p &\equiv (q \rightarrow p) \vee (r \rightarrow p)
 \end{aligned}$$

⊛ $(p \rightarrow q) \equiv (\neg p) \vee q$ disjunctive version of implies

⊛ $(p \rightarrow q) \equiv (\neg q \rightarrow \neg p)$

⊛ $\neg(p \rightarrow q) \equiv \neg(\neg p \vee q) \equiv p \wedge \neg q$

logical equivalences $\square \equiv \triangle$

① show logically equiv.

a) truth table for $\square \leftrightarrow \triangle$ is tautology

b) ask when is \square true? } same conditions?
 when is \triangle true? } \rightarrow logically equiv.

or
 ask when false?

$$\neg (p \rightarrow q) \equiv p \wedge \neg q$$

Wkt.:

p	q	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

$\neg (p \rightarrow q)$
F
T
F
F

when true?

(1) $\neg (p \rightarrow q)$

T when $(p \rightarrow q)$ is false

when p is T and q is F

(2) $p \wedge \neg q$

T when p is T, $\neg q$ is T

when p is T and q is F

(vs) when false?

(1) $\neg (p \rightarrow q)$, F when $p \rightarrow q$ is true

when (i) $p \equiv T \wedge q \equiv T$

(ii) $p \equiv F \wedge q \equiv T$

(iii) $p \equiv F \wedge q \equiv F$

(or)

$$(S \rightarrow (\neg r \wedge \neg b)) \wedge \neg ((\neg r \wedge \neg b) \rightarrow S)$$

$$\neg (p \rightarrow q) \equiv p \wedge \neg q$$

$$\equiv (S \rightarrow (\neg r \wedge \neg b)) \wedge ((\neg r \wedge \neg b) \wedge \neg S)$$

$$\equiv (\neg S \vee (\neg r \wedge \neg b)) \wedge \neg r \wedge \neg b \wedge \neg S$$

$$\equiv (\neg S \vee \neg r) \wedge (\neg S \vee \neg b) \wedge \neg r \wedge \neg b \wedge \neg S$$

\equiv

turning prop function into a proposition

(1st) $\text{Safe}(x)$: "x is safe" U.D of x is actors.

(2nd) evaluator: $\text{Safe}(\text{doug math})$: "doug math is safe"

Quantification [(2nd) all x of U.D. has predicate P (Universal)
(3rd) some x of U.D. has predicate P (Existential)