

# Math 321

Q5

4.3 (33d)

$$\gcd(12345, 54321)$$

$$\gcd(12345, 54321)$$

$$54321 = 9 \cdot 12345 + 6$$

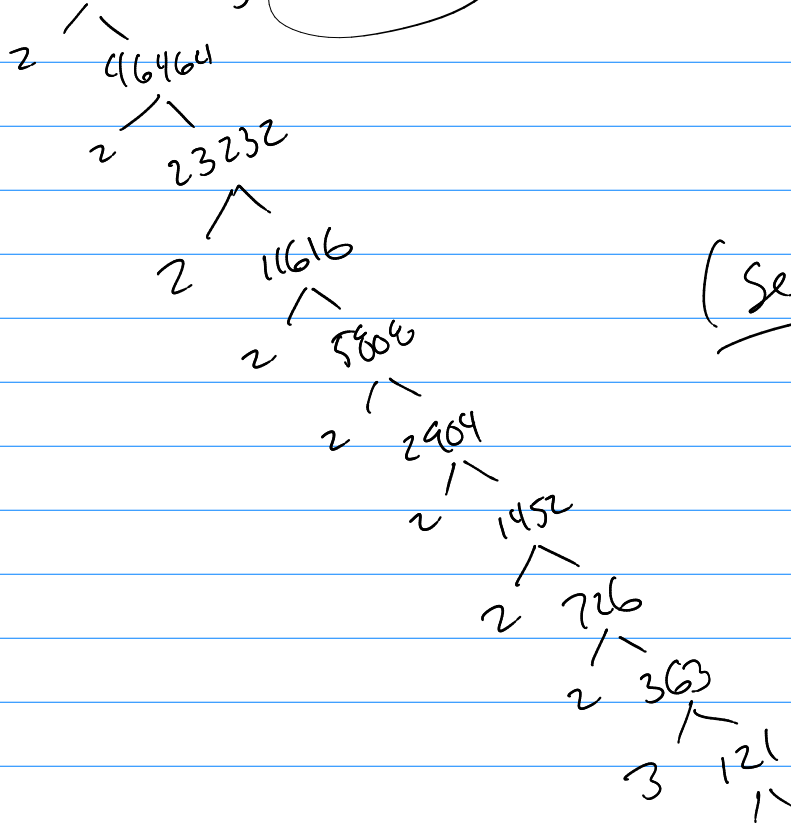
$$12345 = 7 \cdot 6 + 3$$

(see video)

4.3 #29

Prime factors

$$92128, 123552 \rightarrow 2^8 \cdot 3^1 \cdot 11^2 \cdot 13^0, 2^5 \cdot 3^3 \cdot 11^1 \cdot 13^1$$



(see video)

(3e)  $\gcd(117, 213) = 3$

$$3 = s \cdot 213 + t \cdot 117$$

$$213 = 1 \cdot 117 + 96$$

$$117 = 1 \cdot 96 + 21$$

$$96 = 4 \cdot 21 + 12$$

$$21 = 1 \cdot 12 + 9$$

$$12 = 1 \cdot 9 + 3$$

$$9 = 3 \cdot 3 + 0$$

$$3 = 12 - (21 - 12)$$

$$3 = 2 \cdot 12 - 1 \cdot 21$$

$$3 = 2(96 - 4 \cdot 21) - 1 \cdot 21$$

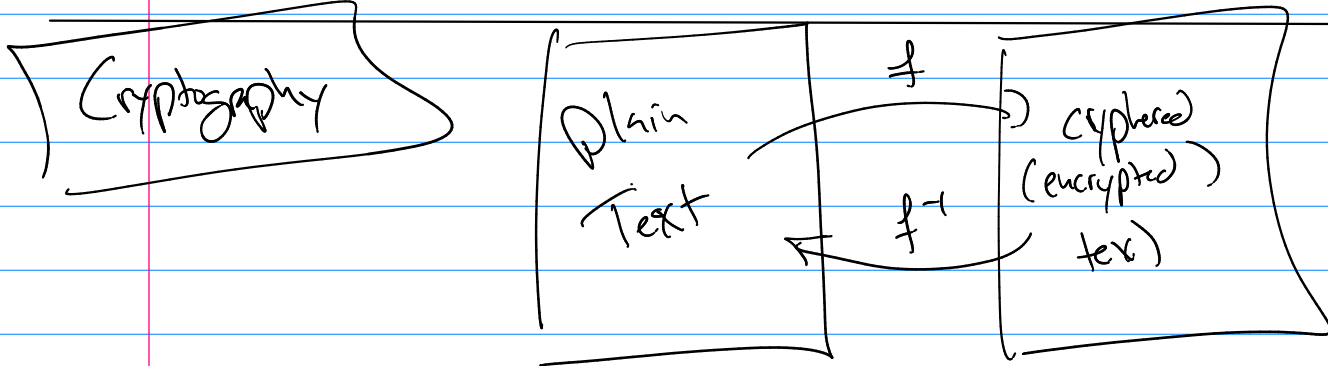
$$3 = 2 \cdot 96 - 9 \cdot 21$$

$$3 = 2 \cdot 96 - 9(117 - 96)$$

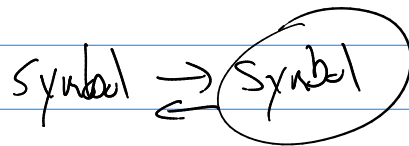
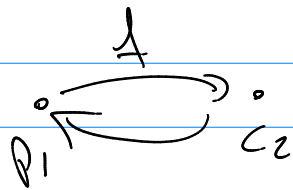
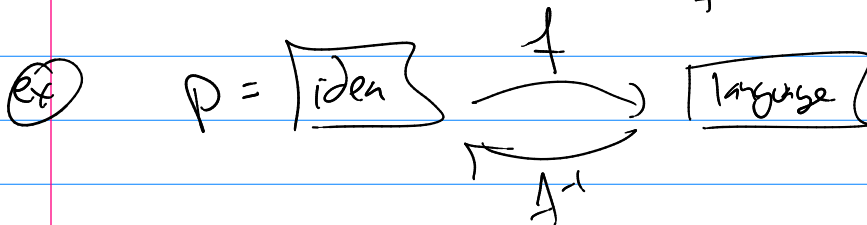
$$3 = 11 \cdot 96 - 9 \cdot 117$$

$$3 = 11 \cdot (213 - 117) - 9 \cdot 117$$

$$3 = 11 \cdot 213 - 20 \cdot 117$$



$$f^{-1}(c) = p \cdot \begin{matrix} \xrightarrow{f} \\ \xleftarrow{f^{-1}} \end{matrix} c = f(p)$$



$$p_i, c_i \in \text{Alphabet}$$

$$\text{Alphabet} = \{s_1, s_2, s_3, \dots, s_n\} \quad |\text{Alphabet}| = n$$

ex) Alph =  $\{a, b, c, \dots, z, \omega, \cdot, \circ, ', \", !, ?\}$   
 $\{A, B, C, \dots, Z\}$

or use table to replace alphabet ex)  $a, b, c, d, \dots, z$   
 $0, 1, 2, 3, \dots, 25$

# Private key cryptography

(see video)

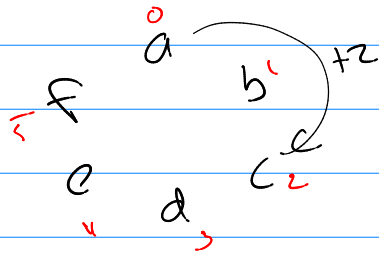
$n = |\text{alphabet}|$

① Shift Cypher

$$c = (p + k) \bmod n$$

$$p = (c - k) \bmod n$$

(ex)



$$c = (p + z) \bmod 26$$

$$p = (c - z) \bmod 26$$

② Affine Shift:

$$f(p) = (ap + k) \bmod n$$

$$f^{-1}(c) = (a^{-1}(c - k)) \bmod n$$

$a^{-1}$  is  $a$ 's inv. mod  $n$

(ex)

$$f(p) = (3 \cdot p + 7) \bmod 11$$

$$f^{-1}(c) = (4 \cdot (c - 7)) \bmod 11 = (4(c - 7)) \bmod 11$$

3's inv. mod 11

$$3 \cdot (?) \equiv 1 \pmod{11}$$

$$11 = 3 \cdot 3 + 2$$

$$1 = 3 - 2$$

$$3 = 1 \cdot 2 + 1$$

$$1 = 1 \cdot 3 - (11 - 3 \cdot 3)$$

$$2 = 2 \cdot 1 + 0$$

$$1 = 4 \cdot 3 - 1 \cdot 11$$

③ random replacement:

(permutation)

$$a \leftrightarrow g$$

$$b \leftrightarrow m$$

$$c \leftrightarrow c$$

⋮

$$z \leftrightarrow y$$

weak to cryptanalysis

④ One-time pad

$$\text{message} = p_1 p_2 p_3 \dots p_k$$

$$\text{random \#s} = r_1 r_2 r_3 \dots r_k$$

$$c_1 = (p_1 + r_1) \bmod n$$

$$c_2 = (p_2 + r_2) \bmod n$$

⋮

$$c_k = (p_k + r_k) \bmod n$$

---