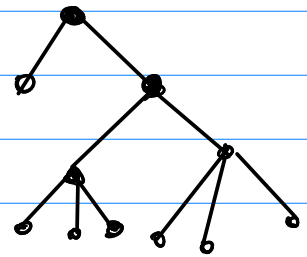
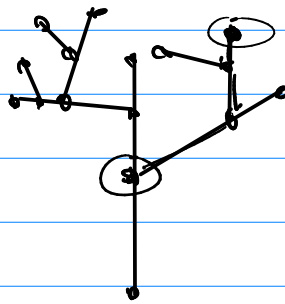
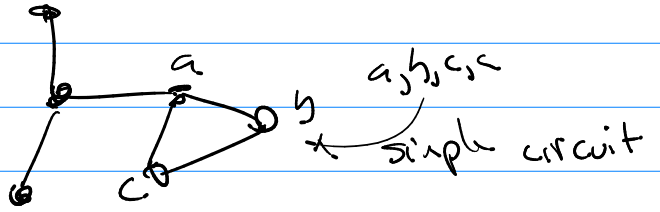
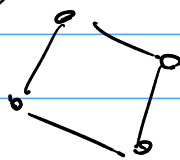


Math 322

$G = (V, E)$
 Ch 11 trees



not trees



Def: tree is a connected undirected graph with no simple circuits

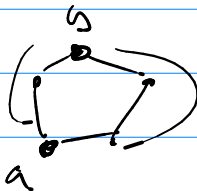
Thm: G is a tree iff there is a unique simple path between any two of its vertices.

Same \Rightarrow .. (let's assume connected undirected graphs)

"no simple circuits iff unique simple path between any two vertices"

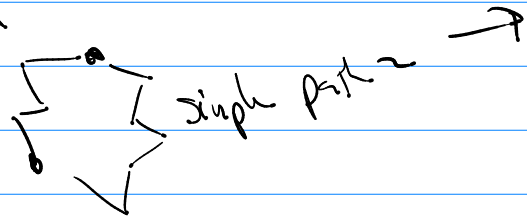
Same \Rightarrow .. "simple circuits iff 2 or more simple paths between two vertices"

PF case "left \Rightarrow right"

assume simple circuit  show 2 or more simple paths (take left or right)

"right → left"

graph part 1



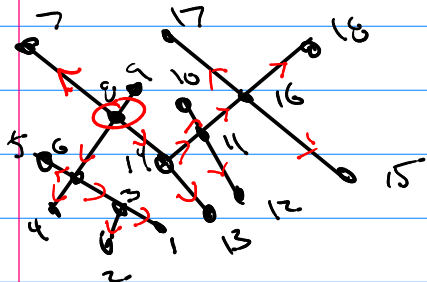
it is a simple circuit



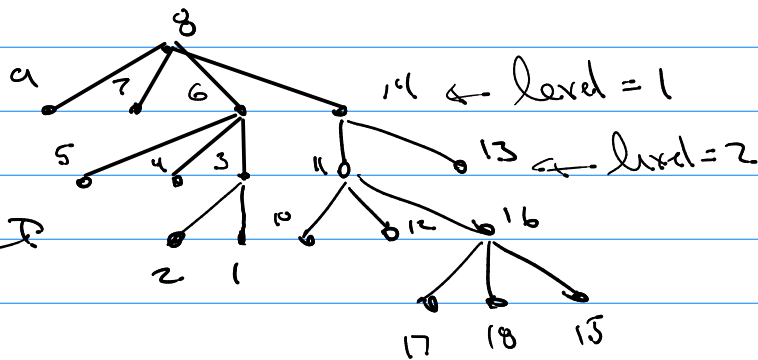
Types of trees

undirected

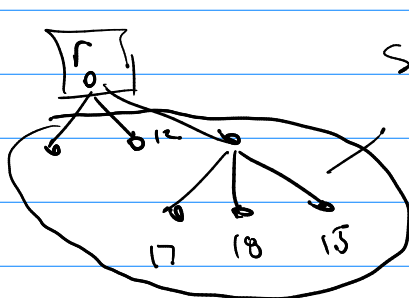
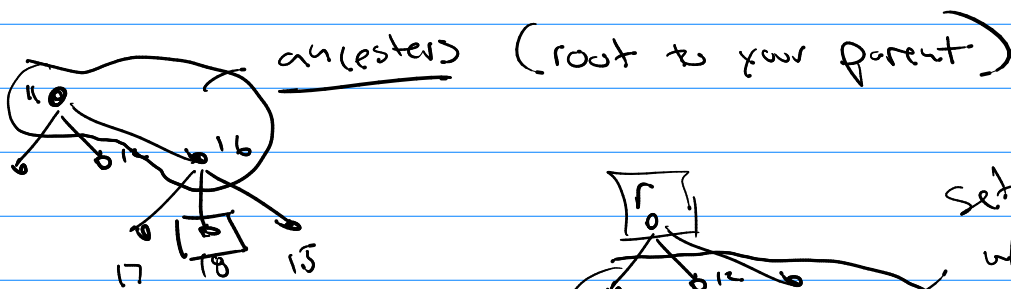
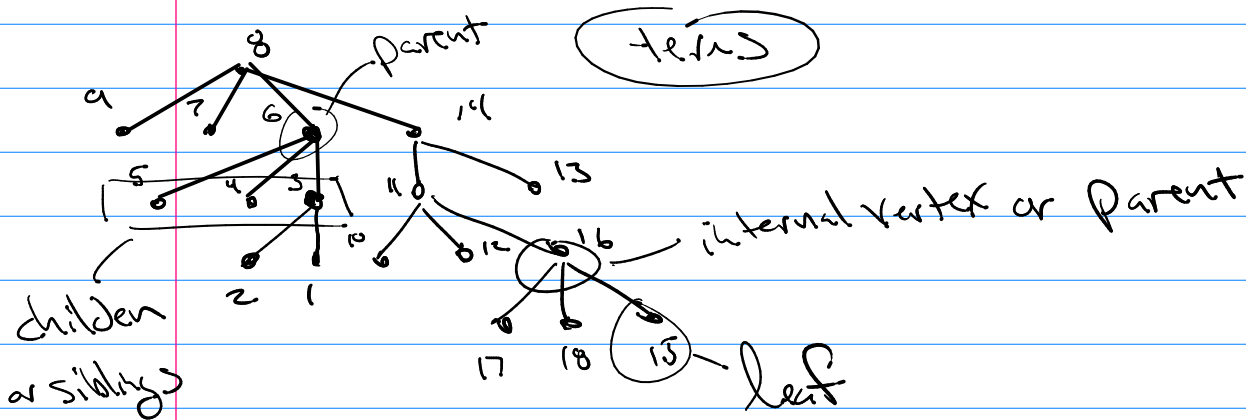
Def A rooted tree is a tree where you call a vertex a root and direct every edge away from the root.



$h=4$



terms



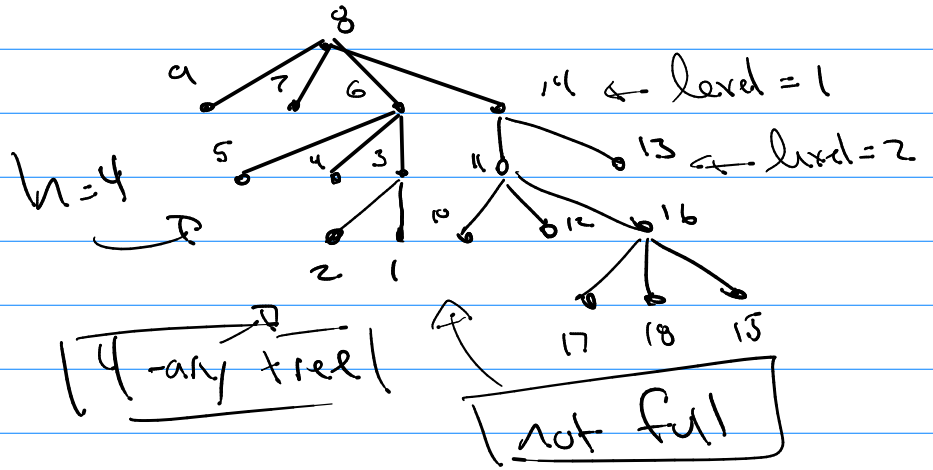
Set of all vertices w/ r as an ancestor: descendants

Key terms (1) level = length of orig. path to root.

(2) height = h = max level

(3) **M-ary tree**

m = largest group of children

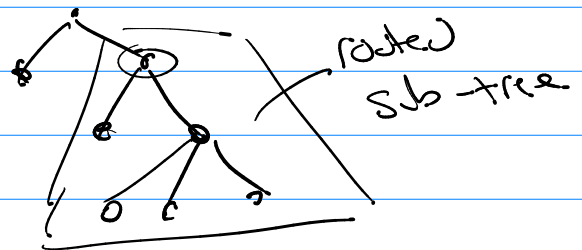
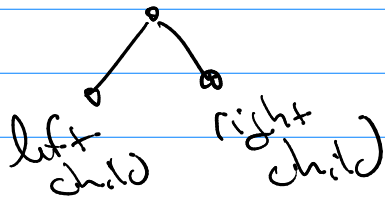


(4) **full m-ary tree**

all parents have m children

Def

ordered rooted tree. Rooted tree, but children are ordered.



Def: rooted sub-trees

Sub-graph taking a vertex as a root and all its descendants.

Properties of trees (\mathcal{H}_2^1) $T = (V, E)$

① $n = |V| \quad n = i + l$
 ↑ ↑
 internal leaves

② \mathcal{H}_2^1 $T = (V, E)$ is a tree with $|V| = n$

a) |children| = $n - 1$ ^{← root}

b) $|E| = n - 1$

③ \mathcal{H}_m^1 T is a full m-ary tree with i internal vertices

a) |children| = $i \cdot m$

b) $n = mi + 1$

\mathcal{H}_m^1 #4 know: $\begin{cases} n = i + l \\ n = mi + 1 \end{cases}$ n, i, l are unknown
 (full m-ary tree)

(i) given $n \rightarrow$ solve for i, l

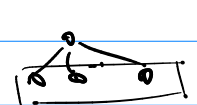
(ii) given $i \rightarrow$ solve for n, l

(iii) given $l \rightarrow$ solve for n, i

⑤ \mathcal{H}_m^1 T is of height h , there is at most m^h leaves

$$\left[\begin{array}{l} l \leq m^h \\ h = 0, 1, 2, \dots \end{array} \right]$$

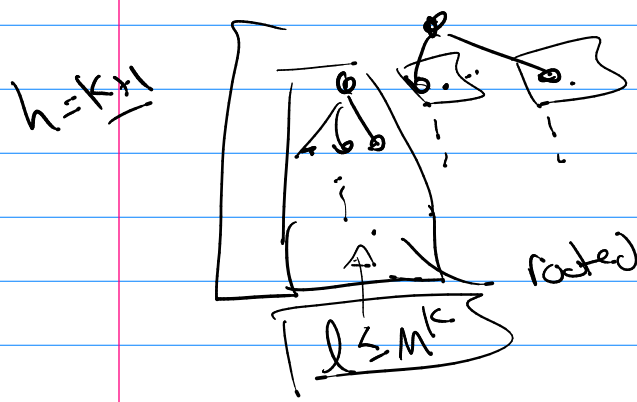
PF Inds: $(h=0) \rightarrow T$ is \bullet $??$ $1 \leq m^0$

base $(h=1) \rightarrow T$ is  $??$, $l \leq m^1$ true!
 children $\leq m$ true!

Inductive

assum: $h=K$ $(l \leq M^K)$ (tree of height K)

show for tree of $h=K+1$



rooted subtrees of 1^{st} children are $h=K$

at most $M \cdot M^K$ leaves M^{K+1}

Def balanced tree has all leaves at level h or $h-1$.

so

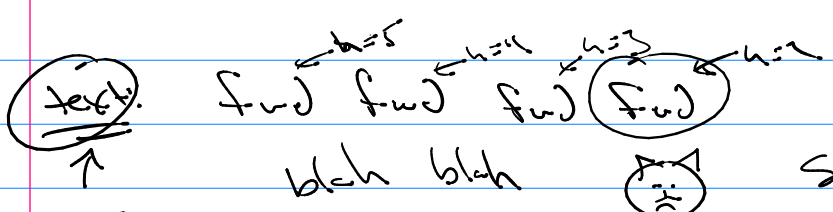
$$l \leq M^h$$

$$\log_M l \leq h$$

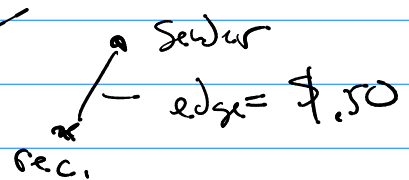
Corollary

$$\text{so } h \geq \lceil \log_M l \rceil$$

$$\text{full and balanced } h = \lceil \log_M l \rceil$$



\$1.25



send to 10 people!

full 10-way tree

how much money = $\boxed{|\mathcal{E}|} (\$.50)$

$$|\mathcal{E}| = \mathcal{N} - 1$$

↑ number of
people & tree

assum. full 10-ary tree & $h = 5$ money = $50k(n-1)$

$$\boxed{l \leq 10^5}$$

↪ worst case assumption $l = 100,000$ (get, not send)

$$\begin{cases} n = i + 10^5 \\ n = 10i + 1 \end{cases}$$

$$9i = 10^5 - 1$$

$$i = \frac{10^5 - 1}{9} = \frac{99,999}{9} = \boxed{11,111} \text{ sends.}$$

$$n = 111,111 \quad \$ = 50k(111,110)$$
