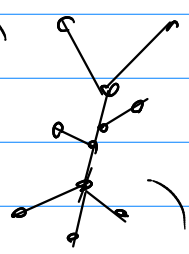
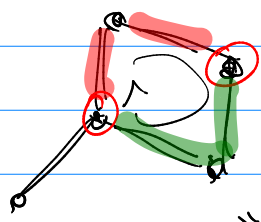


Math 322

Trees

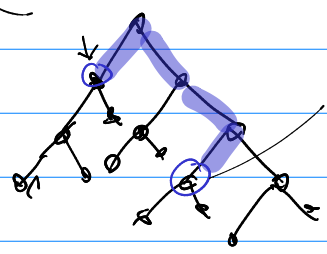
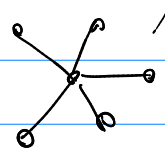


"tree"



simple circuit.

not a "tree"



Def: G is a tree if it is a connected undirected graph with no simple circuits.

Thm

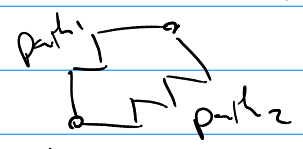
G is a tree if and only if there is a unique simple path between any two vertices.

IPF

Case 1: tree \rightarrow unig. simple path

contradiction \neg unig $\rightarrow \neg$ tree

if \neg unig. simple path says between some two vertices there are 2 or more simple paths.



when you follow both paths you have a simple circuit $\rightarrow \neg$ tree.

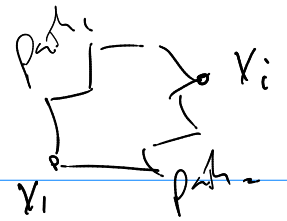
Case 2

(unig simple path \rightarrow tree)

contradiction \rightarrow

\neg tree $\rightarrow \neg$ unig. simple path

tree says we have a simple circuit

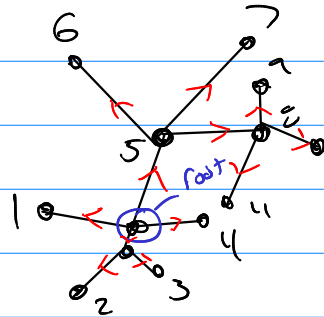


that has another vertex in it.

implies non-uniq simple paths between v_1 and v_i .

Rooted Tree

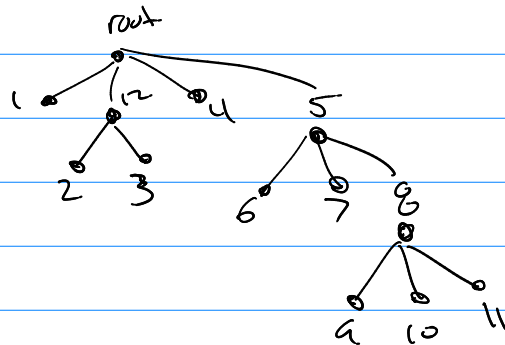
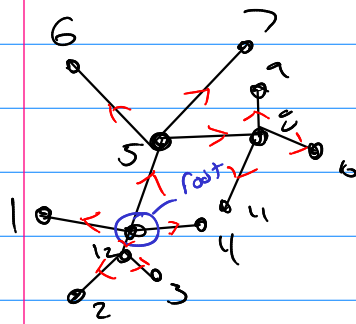
tree (undirected)



① to make a rooted tree
pick one vertex as "root"

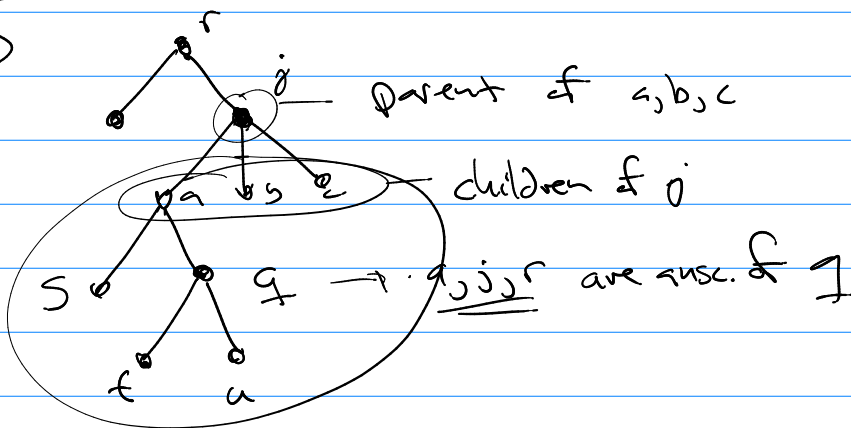
② b/c unig. simple path between any two vertices
→ unig path from root to any vertex.

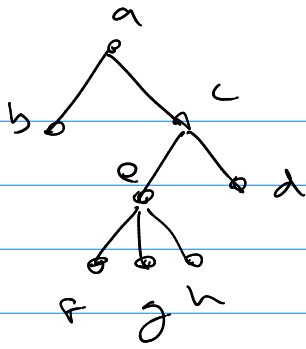
③ direction is 'away' from root.



direction
is down
↓

Terms





internal vertex = have at least one child

(ex) $i = \{a, c, e\}$

leaf = have no children

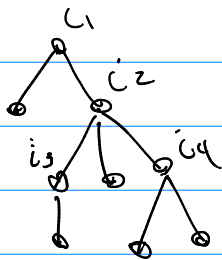
(ex) $l = \{b, d, f, g, h\}$

total vertices $n = i + l$

M-ary rooted tree

→ each internal vertex has at most M children

(ex)



3-ary tree

$i = 4$

$l = 5$

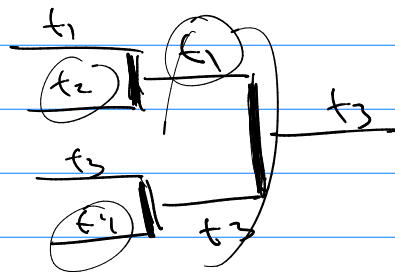
$n = 9$

Note: each parent has a uniq. simple path (single edge) to their child.

So $(n-1)$ children (all - root) ^{not a child}

$|E| = n - 1$

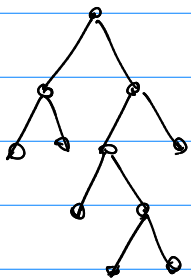
(ex)



Full m-ary rooted tree

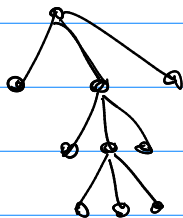
each internal vertex has exactly m children.

(ex) full 2-ary (binary) rooted tree



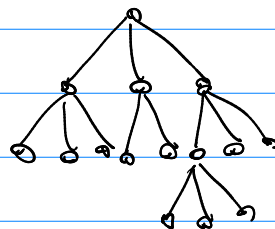
Balanced m-ary tree

all leaves are at bottom or one above bottom.



not balanced

full 3-ary



balanced

3-ary

def.

level = 1

level = 2

level = 3

height = 3

Useful th^m

$|V| = n$

$|E| = n - 1$

$n = i + l$

th^m

full m -ary tree with i internal vertices

$$n = (m \cdot i + 1)$$

children rest

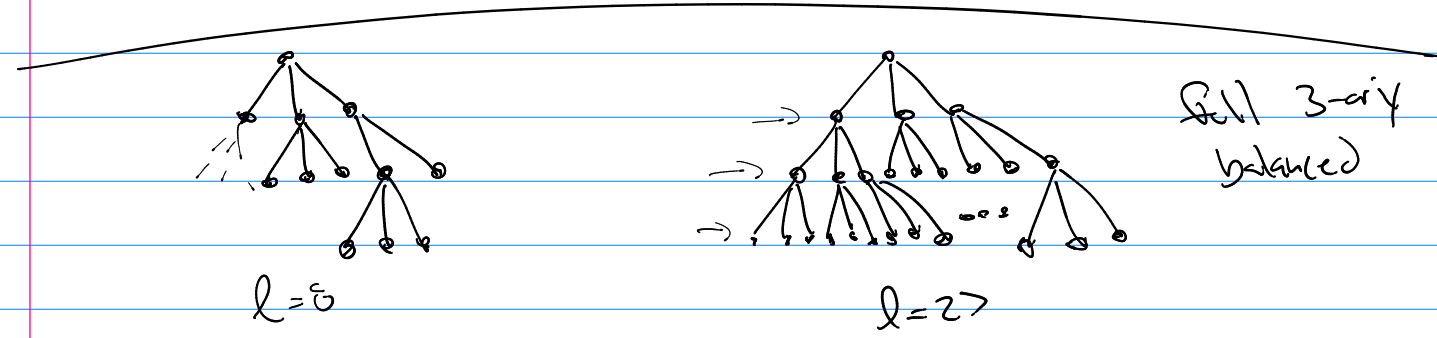
thⁿ #4 p. 753 full M-ary tree with

- ① n vertices, then $i = \frac{n-1}{M}$, $l = \frac{(M-1)n+1}{M}$
- ② i internal, then $n = Mi+1$, $l = (M-1)i+1$
- ③

$n = i+l$ $n = Mi+1$ Know these two things

ex) full 7-ary tree $n = i+l$
 $n = 7i+1$

① $n = 70$ $\begin{cases} 70 = i+l \\ 70 = 7i+1 \end{cases}$

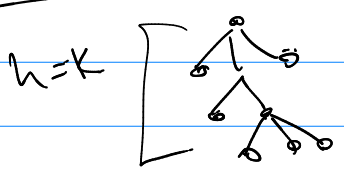


$l \leq M^h$, $h = \text{height of tree}$ ($h = 0, 1, 2, 3, \dots$)

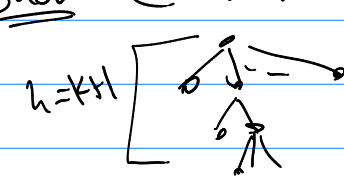
DF (induction)

Base: $\{h=0\}$ tree: $\begin{matrix} \text{root} \\ \bullet \end{matrix}$ $l = 1$
 $M^0 = 1$
 $l \leq M^0 \rightarrow 1 \leq 1$ true

Inductive: assume @ $h=k$ $l \leq M^k$



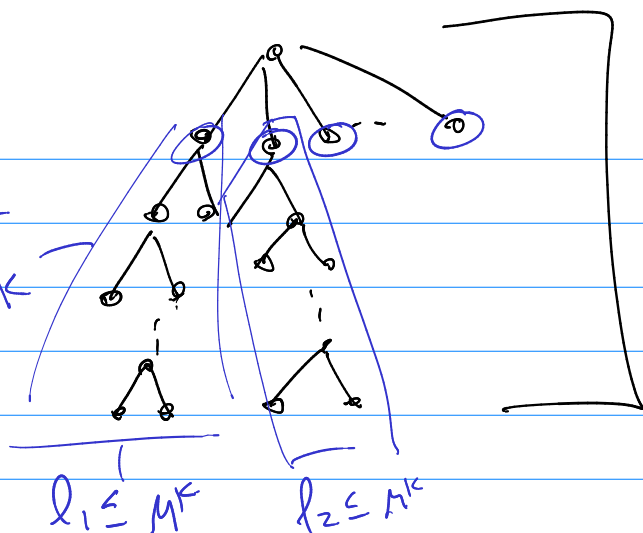
Show @ $h=k+1$



$l \leq M^{(k+1)}$

tree of $h = k+1$

tree of at most height k



$h = k+1$

height $k+1$ has at most $l \leq l_1 + l_2 + \dots + l_n$
 $l \leq M^k + M^k + \dots + M^k$
 $l \leq M^{k+1}$

$$l \leq M^h \Rightarrow \log_M l \leq h$$

$h \geq \lceil \log_M l \rceil$

\Rightarrow of full and balanced $h = \lceil \log_M l \rceil$

ex text:

Fwd; Fwd; Fwd; Fwd



send to 10 people or ...

$h=5$



$n=10$

full 10-ary.

$l \leq M^h$

$l \leq 10^5 = 100000$

$n = (i) + (l)$

sent text

do not send text

$|E| = n-1$

all children

worst case \rightarrow balanced $l = 100,000$

$$n = 10 \quad h = 5 \quad l = 100,000 \quad \left. \vphantom{\begin{matrix} n = 10 \\ h = 5 \\ l = 100,000 \end{matrix}} \right\} \begin{matrix} n = 7 \\ i = 6 \end{matrix}$$

$$\left. \begin{matrix} n = i + l \\ n = ni + 1 \end{matrix} \right\} \rightarrow \begin{cases} n = i + 100,000 \\ n = 10i + 1 \end{cases}$$

$$0 = 9i - 99,999$$

$$i = 11,111$$

$$n = 111,111$$

$$c = 11,111$$

$$l = 100,000$$

edge = text $\$0.10$ per person \rightarrow $\$0.120$ / text

(edges) = 111,110 texts go at

Cost $\$2,322$

