

Math 322

Boolean Algebra:

objects: set $B = \{0, 1\}$ bits = $\{0, 1\}$

rules: a) two binary operators \wedge, \vee object \wedge object
 b) one unary operator $\overline{\quad}$ object \vee object.

boolean variable x is 0 or 1

typically

define ops

x	x/y	$x \wedge y$	$x \vee y$	\overline{x}
1	1	1	1	0
0	0	0	0	1

check laws

$$\overline{\overline{x}} = x$$

double negation law

x	\overline{x}	$\overline{(\overline{x})}$
1	0	1
0	1	0

equal

$$x \vee x = x \quad x \wedge x = x \quad \text{idempotent}$$

$$x \vee \overline{x} = 1$$

$$x \wedge \overline{x} = 0$$

x	\overline{x}	$x \vee \overline{x}$
1	0	1
0	1	1

$x \wedge y = \overline{\overline{x \vee y}}$ DeMorgan's

x	y	$x \wedge y$	$\overline{\overline{x \wedge y}}$	\overline{x}	\overline{y}	$\overline{\overline{\overline{x} \vee \overline{y}}}$
1	1	1	1	0	0	0
0	1	0	0	1	0	0
1	0	0	0	0	1	0
0	0	0	0	1	1	0

equal

Minimal facts to have a Boolean Algebra.

objects: $B = \{0, 1\}$

ops: $\wedge, \vee, \bar{}$

and the following hold:

$$\left. \begin{array}{l} (1) \quad b \vee 0 = b \\ \quad \quad b \wedge 1 = b \end{array} \right\} \text{Identity laws}$$

$$\left. \begin{array}{l} (2) \quad b \vee \bar{b} = 1 \\ \quad \quad b \wedge \bar{b} = 0 \end{array} \right\} \text{Complement laws}$$

$$\left. \begin{array}{l} (3) \quad (b_1 \wedge b_2) \wedge b_3 = b_1 \wedge (b_2 \wedge b_3) \\ \quad \quad (b_1 \vee b_2) \vee b_3 = b_1 \vee (b_2 \vee b_3) \end{array} \right\} \text{Associative laws}$$

$$\left. \begin{array}{l} (4) \quad b_1 \wedge b_2 = b_2 \wedge b_1 \\ \quad \quad b_1 \vee b_2 = b_2 \vee b_1 \end{array} \right\} \text{Commutative laws}$$

$$(5) \text{ distributive laws} \quad (b_1 \wedge (b_2 \vee b_3)) = (b_1 \wedge b_2) \vee (b_1 \wedge b_3)$$

$$(b_1 \vee (b_2 \wedge b_3)) = (b_1 \vee b_2) \wedge (b_1 \vee b_3)$$

So $b \wedge b = b$ (idempotent) could be shown by only using above 5.

$$\begin{aligned} \text{Start with } b &= b \wedge 1 \quad (\text{by identity}) \\ &= b \wedge (b \vee \bar{b}) \quad (\text{by comp. law}) \\ &= (b \wedge b) \vee (b \wedge \bar{b}) \quad (\text{by distrib.}) \\ &= (b \wedge b) \vee 0 \quad (\text{by comp. law}) \\ &= b \wedge b \quad (\text{by identity}) \end{aligned}$$

back to tyrael... (digital design)

objects = 0, 1
ops = +, ·, $\bar{\quad}$

bit-table

x	y	\bar{x}	$x+y$	$x \cdot y$
1	1	0	1	1
1	0	0	1	0
0	1	1	1	0
0	0	1	0	0

use bit tables to verify all the laws.

Boolean Expression: combination of bits, variables, ops.

(ex) ① $x + [(y \cdot z) \cdot 1] + \bar{0}$

② $x + \bar{y}$

③ $(\overline{x+y}) \cdot (\bar{x} + 0)$

bit table:

x	y	\bar{x}	0	$\bar{x} + 0$	$x + y$	$\overline{x+y}$	$(\overline{x+y}) \cdot (\bar{x} + 0)$
1	1	0	0	0	1	0	0
1	0	0	0	0	1	0	0
0	1	1	0	1	1	0	0
0	0	1	0	1	0	1	1

(x,y) expression

$(1,1) \rightarrow (0)$

$(1,0) \rightarrow (0)$

$(0,1) \rightarrow (0)$

$(0,0) \rightarrow (1)$

function taking (x,y) ordered pairs and return a bit.

Boolean Function: $B^n : (n\text{-tuple of } n\text{-bits}) \rightarrow (\text{bit})$

$f(b_1, b_2, \dots, b_n) =$ expression of b_i boolean variables
boolean function of degree n .

(ex) $f(x, y, z) = \bar{x} + (y \cdot \bar{z}) + \bar{y}$

x	y	z	\bar{x}	\bar{y}	\bar{z}	$y \cdot \bar{z}$	$(y \cdot \bar{z}) + \bar{y}$	$\bar{x} + (y \cdot \bar{z}) + \bar{y}$
1	1	1	0	0	0	0	0	0
1	1	0	0	0	1	1	1	1
1	0	1	0	1	0	0	1	1
1	0	0	0	1	1	0	1	1
0	1	1	1	0	0	0	0	1
0	1	0	1	0	1	1	1	1
0	0	1	1	1	0	0	1	1
0	0	0	1	1	1	0	1	1

$(1, 1, 1) \xrightarrow{f} (0)$
 $(1, 1, 0) \rightarrow (1)$
 $(1, 0, 1) \rightarrow (1)$
 \vdots

$f =$ expression (finite ops \pm , finite (countable) variables)

have countably infinite possible expressions.

(ex) $f(x, y)$ has countably infinite possible expressions

$f(x, y) = 1, f(x, y) = (\overline{x + y + 0}), \dots$

but all of them are functions

have a table

x	y	f
1	1	1 or 0
1	0	1 or 0
0	1	1 or 0
0	0	1 or 0

$(\text{tables}) = 2^{|\text{rows}|} = 2^4 = 16 =$

f is B^n (n-tuples input)

$(\text{tables}) = 2^{|\text{rows}|} = 2^{2^{|\text{rows}|}} = 2^{2^n}$ (finite)

ex Program = $f: B^n \rightarrow B$

only 2^{2^n} total "exist",

So... if only 2^{2^n} exist (unique boolean tables)
 Can we find a "best" expression for a given table.

(ex) $f =$ (above expression)

x	y	z	f
1	1	0	0
1	0	0	1
0	1	0	1
0	0	0	0

Min-term / Maxterm expansions

x	y	f
1	0	
0	0	
0	1	
1	1	

each are 0 or 1

Consider $X \cdot y$

x	y	$X \cdot y$	$\overline{X \cdot y}$	$\overline{X} \cdot y$	$X \cdot \overline{y}$
1	1	1	0	0	0
1	0	0	1	0	0
0	1	0	0	1	0
0	0	0	0	0	1

Consider each row has a minterm that is its one

x	y	Minterm of this row
1	1	Xy
1	0	$X\overline{y}$
0	1	$\overline{X}y$
0	0	$\overline{X}\overline{y}$

Def Minterm is the product of all variables as x or \overline{x}

Min term expression

X	y	z	f
1	1	0	1
1	0	0	1
0	1	0	0
0	0	0	0
0	0	1	1

$f = xy\bar{z} + x\bar{y}\bar{z} + \bar{x}\bar{y}z$

1 → $xy\bar{z}$
1 → $x\bar{y}\bar{z}$
1 → $\bar{x}\bar{y}z$