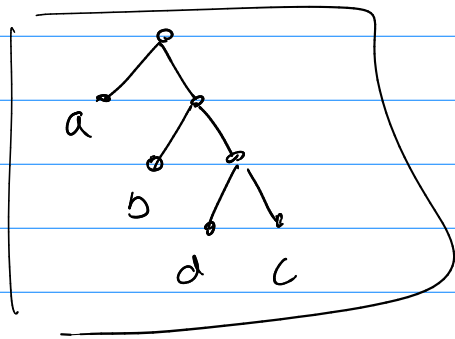
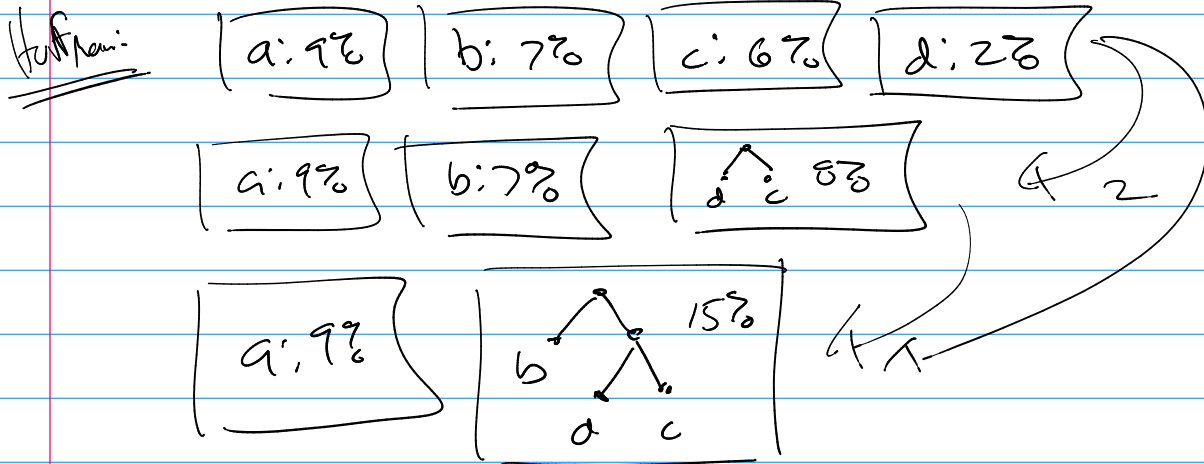


Math 322



17.23 $f: B^n \rightarrow B$ (a) $f(x_1, x_2, \dots, x_n) = \underline{\text{expression}}$

tblw:

	x_1	x_2	...	x_n	f
2^n row	1	1	...	1	0 or 1
	1	1	...	0	
	:	:	:	:	
	0	0	...	0	

total orig. tblw
 $(2)^{2^n}$

So there are only 2^n (orig) $f(x_1, \dots, x_n)$

(e)

x	y	z	$f = (\text{expression})?$
1	1	1	1 ← $x \cdot y \cdot z$
1	1	0	0
1	0	1	1 ← $x \cdot \bar{y} \cdot z$
1	0	0	0
0	1	1	0
0	1	0	0
0	0	1	0
0	0	0	1 ← $\bar{x} \cdot \bar{y} \cdot \bar{z}$

$f(x, y, z) = (x \cdot y \cdot z) + (x \cdot \bar{y} \cdot z) + (\bar{x} \cdot \bar{y} \cdot \bar{z})$

Focus on min-term / max term.

① Min term: product of all variables as literals X or \bar{X}
 ex) $(X_1 \cdot X_2 \cdot \bar{X}_3 \cdot X_4 \cdot \dots \cdot X_n)$
 all X_i as X_i or \bar{X}_i

② Max term: sum of all variables as literals
 ex) $(X_1 + X_2 + \bar{X}_3 + \dots + X_n)$
 all X_i as X_i or \bar{X}_i

Why?

① min terms are only 1 when all literals are 1's
 ex) $(X_1 \cdot X_2 \cdot \bar{X}_3)$ gives 1 only when $X_1 = 1$
 $X_2 = 1$
 $\bar{X}_3 = 1 \rightarrow X_3 = 0$

② max terms are only 0 when all literals are 0's
 ex) $(X_1 + \bar{X}_2 + X_3)$ gives 0 only when $X_1 = 0$
 $\bar{X}_2 = 0 \rightarrow X_2 = 1$
 $X_3 = 0$

<u>So</u>		when 1's	when 0's
	X, y, z	assoc. min term	assoc max term
	1 1 1	$X_1 \cdot X_2 \cdot X_3$	$(\bar{X}_1 + \bar{X}_2 + \bar{X}_3)$
	1 1 0	$X_1 \cdot X_2 \cdot \bar{X}_3$	$(\bar{X}_1 + \bar{X}_2 + X_3)$
	1 0 1	$X_1 \cdot \bar{X}_2 \cdot X_3$	$(\bar{X}_1 + X_2 + \bar{X}_3)$
	1 0 0	$X_1 \cdot \bar{X}_2 \cdot \bar{X}_3$	$(\bar{X}_1 + X_2 + X_3)$
	0 1 1	$\bar{X}_1 \cdot X_2 \cdot X_3$	$(X_1 + \bar{X}_2 + \bar{X}_3)$
	0 1 0	$\bar{X}_1 \cdot X_2 \cdot \bar{X}_3$	$(X_1 + \bar{X}_2 + X_3)$
	0 0 1	$\bar{X}_1 \cdot \bar{X}_2 \cdot X_3$	$(X_1 + X_2 + \bar{X}_3)$
	0 0 0	$\bar{X}_1 \cdot \bar{X}_2 \cdot \bar{X}_3$	$(X_1 + X_2 + X_3)$

A Minterm Expansion / Sum of Products / Disjunctive Normal Form
 → Sum of minterms

(ex)

X	Y	f
1	1	0
1	0	1
0	1	0
0	0	1

$\left. \begin{array}{l} \text{1} \leftarrow x \cdot \bar{y} \\ \text{1} \leftarrow \bar{x} \cdot y \end{array} \right\} f(x,y) = (x \cdot \bar{y}) + (\bar{x} \cdot y)$

X	Y	f
1	1	0
1	0	1
0	1	0
0	0	0

$\text{1} \leftarrow x \cdot \bar{y} \quad f(x,y) = (x \cdot \bar{y})$

Note:
 $\neg(x \rightarrow y)$
 $\rightarrow (\neg x \vee y)$
 $x \wedge \neg y$

B Maxterm expansion / Product of Sums / Conjunctive Normal Form
 → Product Maxterms

X	Y	f
1	1	0
1	0	1
0	1	0
0	0	1

$\left. \begin{array}{l} \text{0} \leftarrow (\bar{x} + \bar{y}) \\ \text{0} \leftarrow (x + y) \end{array} \right\} f(x,y) = (\bar{x} + \bar{y})(x + y)$

X	Y	f
1	1	0
1	0	1
0	1	1
0	0	1

$\text{0} \leftarrow (\bar{x} + y) \quad f(x,y) = (\bar{x} + y)$

(ex) focus on 1's \equiv minterms

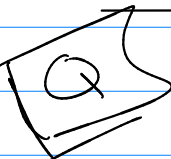
X	Y	Z	f
1	1	1	0
1	1	0	1
1	0	1	0
1	0	0	0
0	1	1	1
0	1	0	0
0	0	1	0
0	0	0	1

$f(x,y,z) = (x \cdot y \cdot \bar{z}) + (\bar{x} \cdot y \cdot z) + (\bar{x} \cdot \bar{y} \cdot \bar{z})$

focus on 0's = use max terms

x	y	z	f
1	1	1	1 (x+y+z)
1	1	0	1 (x+y+z̄)
1	0	1	1 (x̄+y+z)
1	0	0	1 (x̄+y+z̄)
0	1	1	1 (x+y+z)
0	1	0	1 (x+y+z̄)
0	0	1	1 (x̄+y+z)
0	0	0	1 (x̄+y+z̄)

$$f(x,y,z) = (\bar{x}+y+z)(\bar{x}+y+\bar{z})(\bar{x}+y+z)(x+\bar{y}+z)(x+y+\bar{z})$$



How to go from expression into Sum of Minterms? Product of max terms?

ex $f(x,y,z) = (xz) + (y)$

tech #1

use a table

x	y	z	xz	xz+y
1	1	1	1	1
1	1	0	0	1
1	0	1	1	1
1	0	0	0	0
0	1	1	0	1
0	1	0	0	1
0	0	1	0	0
0	0	0	0	0

$$f(x,y,z) = (xyz) + (xy\bar{z}) + (x\bar{y}z) + (\bar{x}yz) + (\bar{x}y\bar{z})$$

$$f(x,y,z) = (\bar{x}+y+z)(x+y+\bar{z})(x+y+z)$$

tech #2

use $b \cdot 1 = b$ $b + 0 = b$
 $\text{and } (b + \bar{b}) = 1$ $(b \cdot \bar{b}) = 0$

ex $f(x,y,z) = (xz) + (y) = (x \cdot 1 \cdot z) + (1 \cdot y \cdot 1)$
 $= (x(y+\bar{y})z) + ((x+\bar{x})y)$
 $= (xyz) + (x\bar{y}z) + (xy\cdot 1) + (\bar{x}y\cdot 1)$
 $= (xyz) + (x\bar{y}z) + (xy(z+\bar{z})) + (\bar{x}y(z+\bar{z}))$
 $= (xyz) + (x\bar{y}z) + (xy\underline{z}) + (xy\bar{z}) + (\bar{x}yz) + (\bar{x}y\bar{z})$
 $= (xyz) + (x\bar{y}z) + (xy\bar{z}) + (\bar{x}yz) + (\bar{x}y\bar{z})$

$$f(x, y, z) = xz + y$$

$$b \wedge 1 = b$$

$$b \vee 0 = b$$

$$b \vee \bar{b} = 1$$

$$b \wedge \bar{b} = 0$$

$$= (x \wedge z) \vee y \quad \text{prod \& Maxterms} \quad \underline{\underline{(x \vee \bar{x}) \wedge (x \vee \bar{x})}}$$

$$= (x \vee y) \wedge (y \vee z)$$

$$= (x \vee y \vee \underbrace{0}_{z \wedge \bar{z}}) \wedge (\underbrace{0}_{x \wedge \bar{x}} \vee y \vee z)$$

$$= (x \vee y \vee (z \wedge \bar{z})) \wedge ((x \wedge \bar{x}) \vee y \vee z)$$

$$= \underline{(x \vee y \vee z)} \wedge (x \vee y \vee \bar{z}) \wedge \underline{(x \vee y \vee z)} \wedge (\bar{x} \vee y \vee z)$$

$$= (x \vee y \vee z) \wedge (x \vee y \vee \bar{z}) \wedge (\bar{x} \vee y \vee z)$$

$$\textcircled{vs} \quad f(x, y, z) = xz + y = (x \cdot z) + y$$

$$= (x+y)(z+y)$$

$$= (x+y+0)(0+y+z)$$

$$= (x+y)(z+\bar{z}) + (x+\bar{x})(y+z)$$

$$= \text{etc}$$

Note: $f: B^n \rightarrow B \quad f(x_1, x_2, \dots, x_n) = \underline{\text{expression}}$

by our work with maxterm / minterm expansions we know that $\forall B \quad 2^n$ unig. functions can be written using

only $x_1, x_2, \dots, x_n, \underbrace{+, \cdot, -}$

\leftarrow call $\{+, \cdot, -\}$ to be functionally complete.

Q can we have fewer?

by DeMorgans: $\overline{X+y} = \overline{X} \cdot \overline{y}$ so $X+y = \overline{\overline{X} \cdot \overline{y}}$

so if you had $\{\circ, \Delta\} \rightarrow f(x,y) = (X+y) = \overline{\overline{X} \cdot \overline{y}}$

b/c $\overline{\square} + \Delta = \overline{\square} \cdot \overline{\Delta}$ $\{\circ, -\}$ is functionally complete

also

by DeMorgans $\overline{X \cdot y} = \overline{X} + \overline{y} \rightarrow X \cdot y = \overline{\overline{X} + \overline{y}}$

b/c $\square \cdot \Delta = \overline{\overline{\square} + \overline{\Delta}}$ $\{+, -\}$ is functionally complete

Carryover

X	y	$X \downarrow y$	$X y$
1	1	0	0
1	0	0	1
0	1	1	0
0	0	1	1

↑ ↑
 not (or) not (And)
 NOR NAND

$$X \downarrow X = \overline{X}$$

$$(X \downarrow y) \downarrow (X \downarrow y) = X + y$$

$$(X \downarrow X) \downarrow (y \downarrow y) = X \cdot y$$

so $\{\downarrow\}$ is functionally complete

$$X | X = \overline{X}$$

$$(X | y) | (X | y) = X \cdot y$$

$$(X | X) | (y | y) = X + y$$

so $\{| \}$ is functionally complete

complete.