

Math 322

Thm Principle of Well ordered Induction

$$\left[\underbrace{P(1^{st})}_{T} \wedge \left\{ \forall k \left(\underbrace{P(1^{st}) \wedge P(2^{nd}) \wedge \dots \wedge P(k^{th})}_{T} \rightarrow P(k+1^{st}) \right) \right\} \right] \rightarrow \forall n P(n)$$

⊗ (use contradiction)

assume: $P(1^{st})$ case is true

assume: Induction Step is true

$$\underbrace{P(1^{st})}_{T} \wedge \underbrace{P(2^{nd})}_{T} \wedge \dots \wedge \underbrace{P(k^{th})}_{T} \rightarrow \underbrace{P(k+1^{st})}_{T}$$

assume: $\forall n P(n)$ is false

$\rightarrow P(n)$ has one or more counter examples

b/c our set is well ordered Our counter examples can be seen as 1st counter example, 2nd c.e., ...

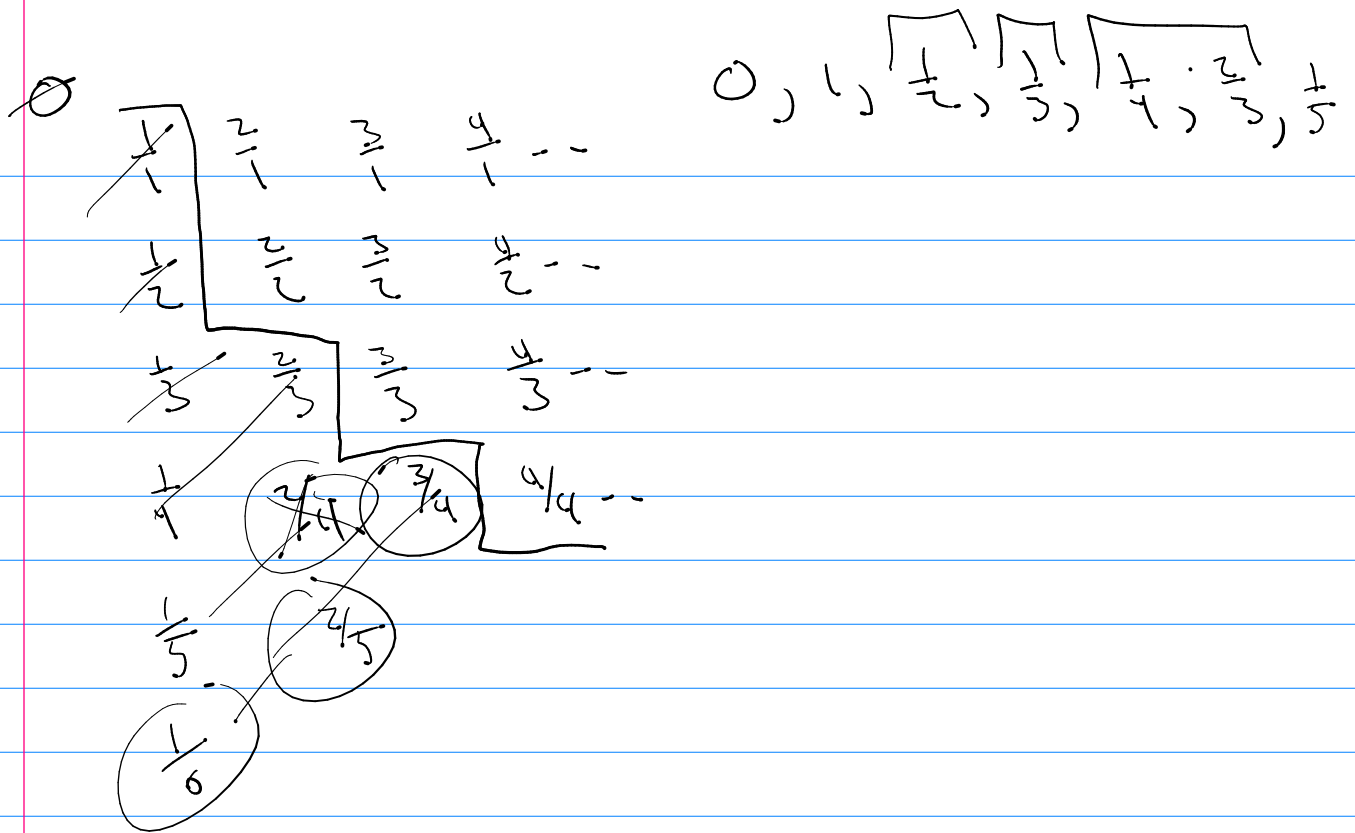
⊗

$$P(n) \quad P(1), P(2), \dots, P(k), \boxed{P(1^{st} \text{ counter ex.})}, \dots$$

$T \quad T \quad \dots \quad T \quad F \quad \nabla$

\swarrow Contradiction

but the inductive step says.. $P(1^{st} \text{ c.e.}) \equiv T$



Lexicographic order (Dictionary Sort)

$n < m$ $a = (a_1, a_2, \dots, a_n)$ $b = (b_1, b_2, \dots, b_m)$

$a < b$

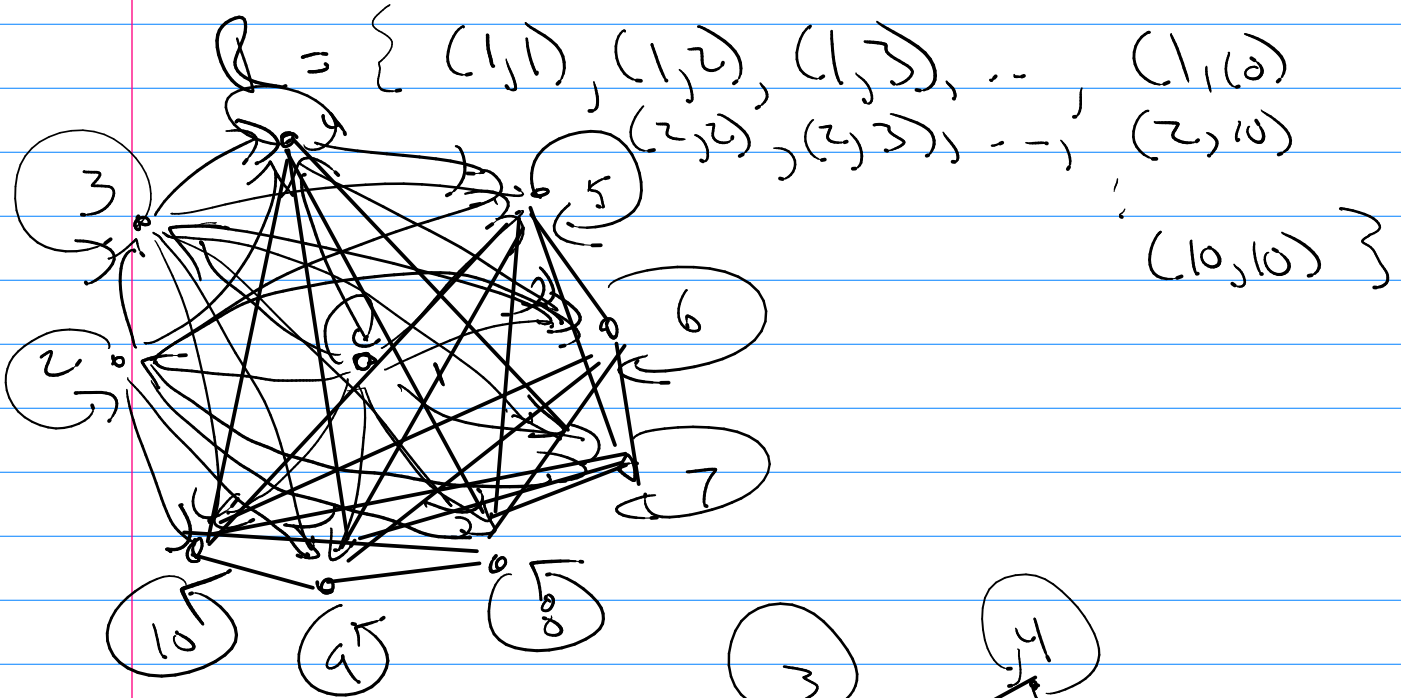
when $a_1 < b_1$
 or if $a_1 = b_1$, $a_2 < b_2$
 or if $a_2 = b_2$, $a_3 < b_3$
 ...

$a_{n-1} = b_{n-1}, a_n < b_n$

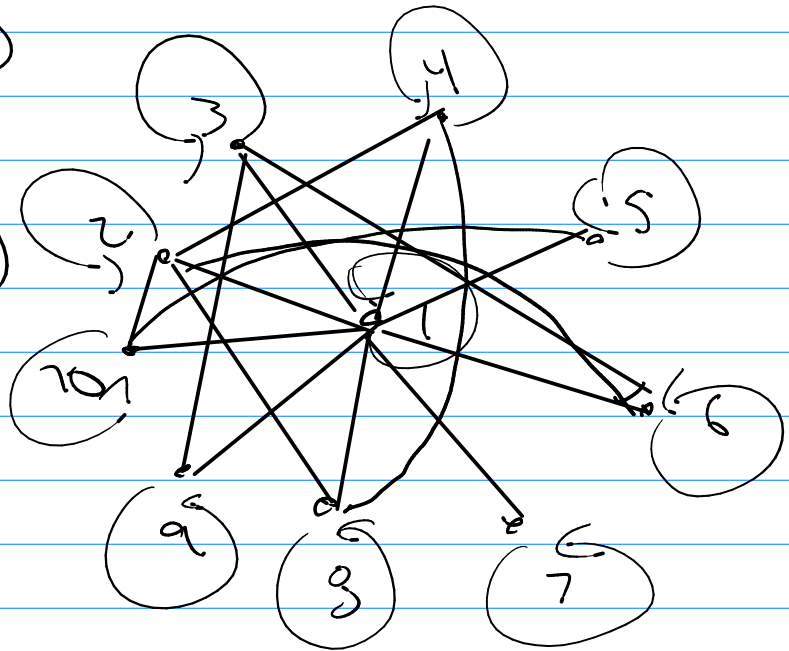
or if $a_n = b_n$ but $M > n$

Hasse Diagrams

$$(\{1, 2, \dots, 10\}, \leq)$$



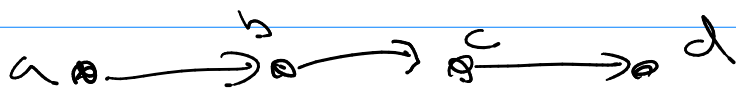
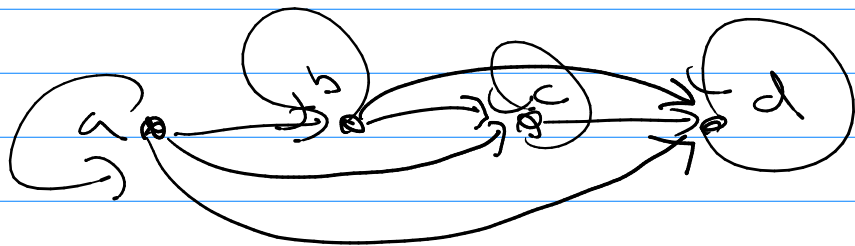
$$(\{1, 2, \dots, 10\}, |)$$



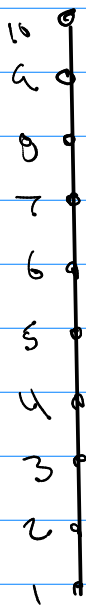
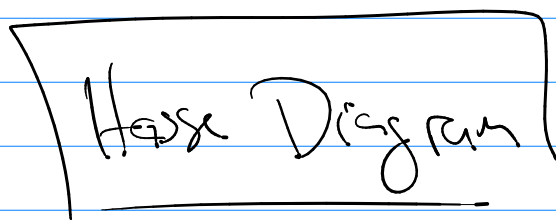
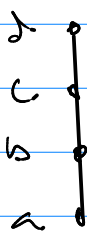
Simplify the graph so the ref, antisym, trans are assumed (b/c (S, \leq)) but not specifically drawn

① ref. (do not draw loops)

② trans $\equiv \forall n \ R^n \subseteq R$

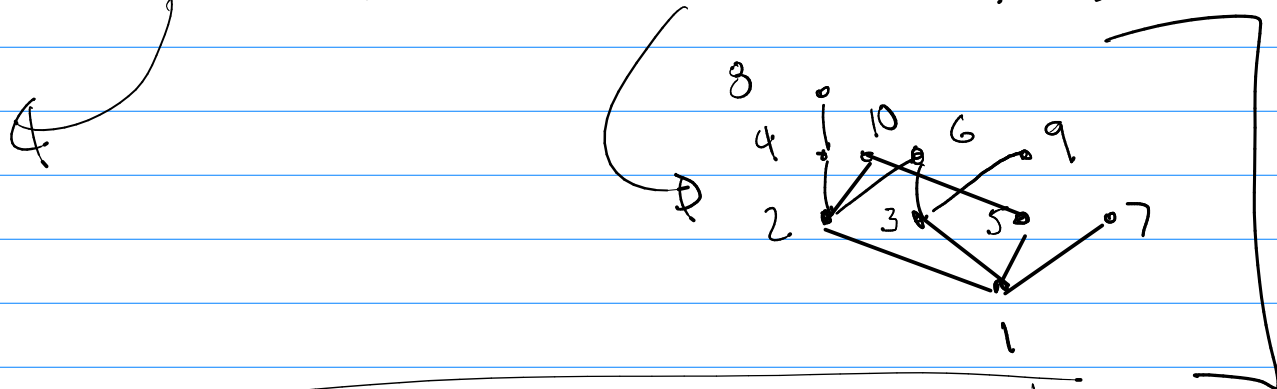


③ antisym. (remove arrows and assume arrows point up)



$(\{1, 2, \dots, 10\}, \leq)$

$(\{1, 2, 3, \dots, 10\}, \mid)$



Extremal Properties of
Hasse Diagrams