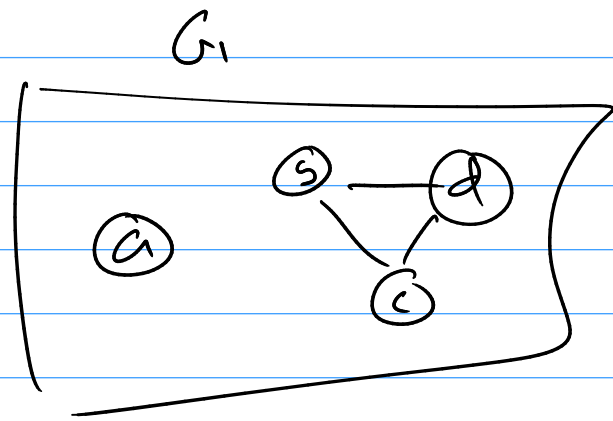
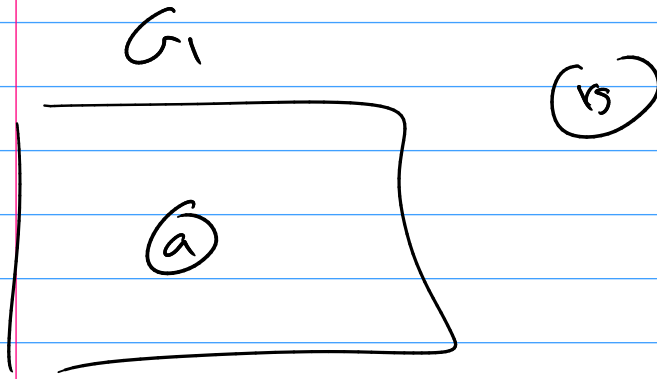


Math 322

Q's

Connected?

Concept is "one piece"



Up Next:

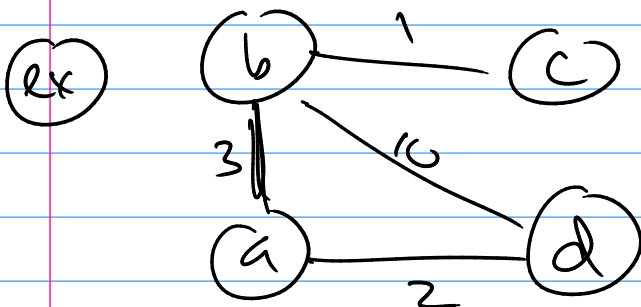
9.5 into ch 10

trees

Weighted graphs:

$$G = (V, E, w(e))$$

↑
weight
function



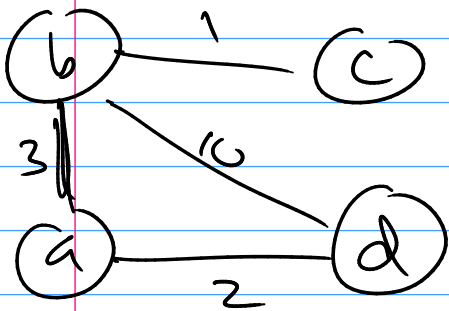
$$w(\{a,b\}) = 3$$

$$w(\{a,d\}) = 2$$

$$w(\{b,c\}) = 1$$

$$w(\{b,a\}) = 3$$

Weighted graph length path is
 Sum of weights in path.



path: a, b, d, b, c, b

before \Rightarrow length = 5 (edges)

Now in a weighted graph

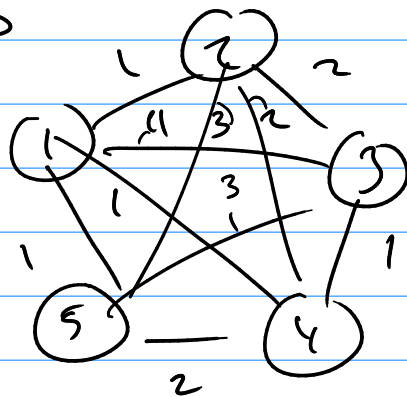
$$\underline{\underline{\text{length}}} = 3 + 10 + 10 + 1 + 1 = 25$$

App Optimization & path lengths

- (1) Minimal paths that Solve problems
- (2) Maximal paths that Solve problems

(ex) traveling salesman problem \rightarrow (n-1)! unig. hamiltonian circuits.
 kn with weights

(ex) 5 city



So 1, 2, 3, 4, 5, 1 length = 7
 or 1, 3, 5, 2, 4, 1 length = 11

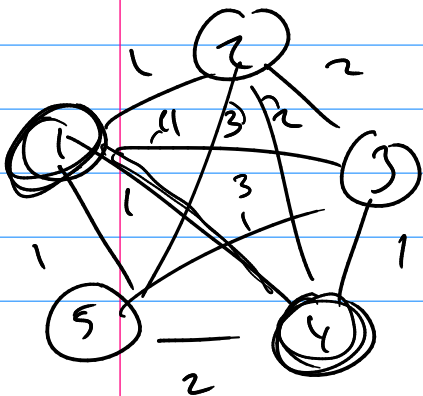
but 4! = 24 Hamiltonian circuits is
 finding $(1 \cdot 3 \cdot 2 \cdot 1) = 24$ circuits
 and their costs. Min = Min
 Max = Max

Approx Ans? Heuristic Algorithms

Analysis of Heuristic Algorithms

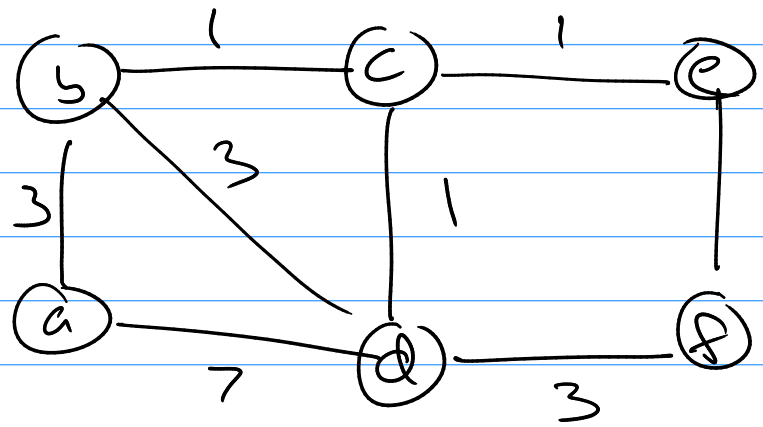
- ① Limit the problem by limiting understanding of $V, E, W(E)$
- ② then approx $\approx \frac{\text{cost}}{1}$ optimal

Ex 3 closest neighbor algorithm



closest neighbor 1, 4, 3, 5, 2, 1
 1, 1, 3, 1
 length = 7

given $G = (V, E, w(e))$ can you find the list of all paths starting at v_i and list each other $v_j \neq v_i$ as a cheapest path?



single.

Start @ (a) → a to a → cheapest paths
 a to b —
 a to c —
 a to d —
 a to e —
 a to f —

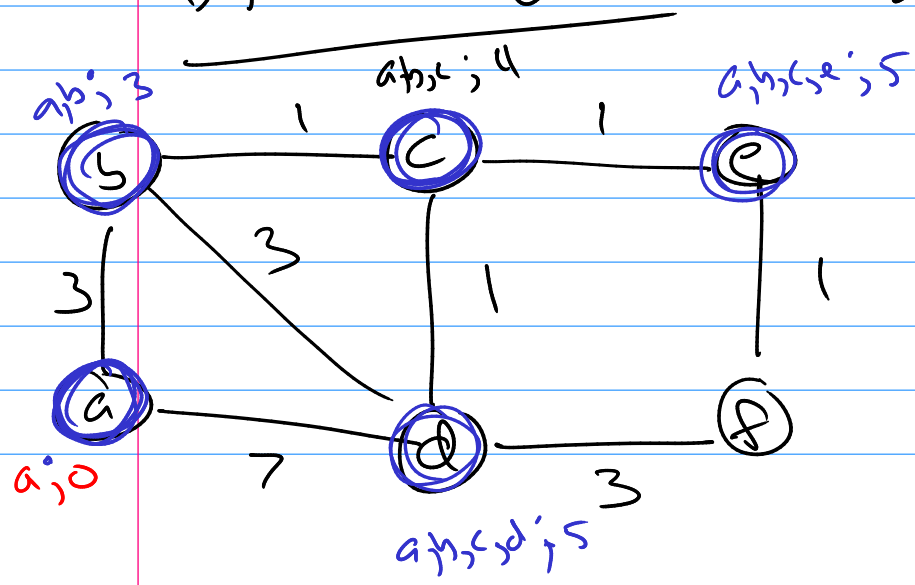
Search

known → unknown
 b,c 4
 b,d 7 a,d 7

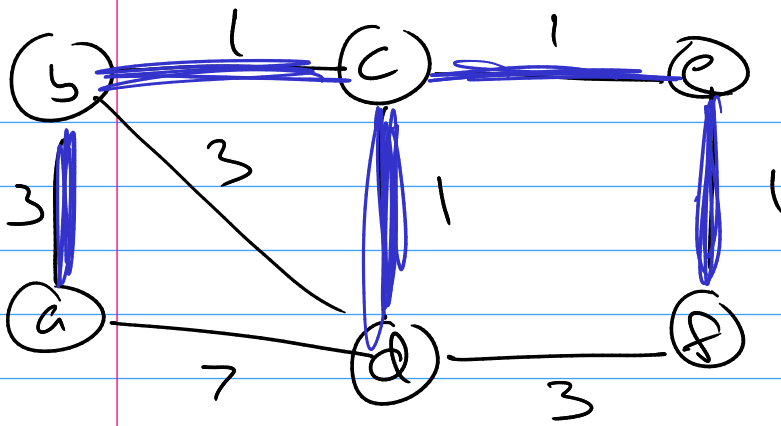
Dijkstra's Algorithm

known
a,b

unknown
c,d,e,f



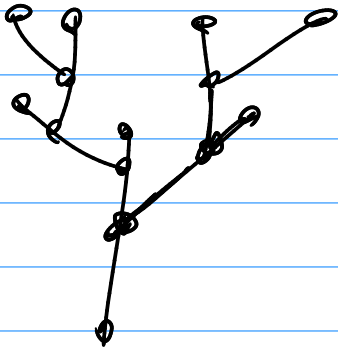
Vertex	Path	Cost
a	a	0
b	a,b	3
c	a,b,c	4
e	a,b,c,e	5
d	a,b,c,d	5
f	a,b,c,d,f	6



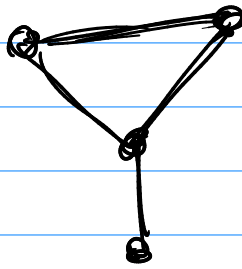
$a: 0$
 $a, b: 3$
 $a, b, c: 4$
 $a, b, c, d, e: 5$
 $a, b, c, d, e, f: 6$

Ch 10 trees

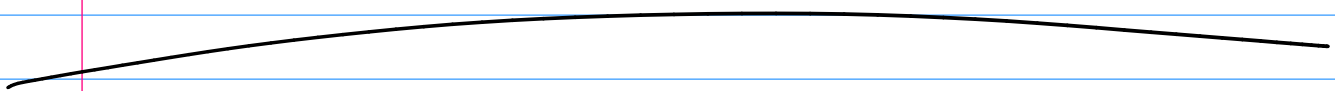
Df: Simple undirected graph with no simple circuits is a tree.



tree



not a tree



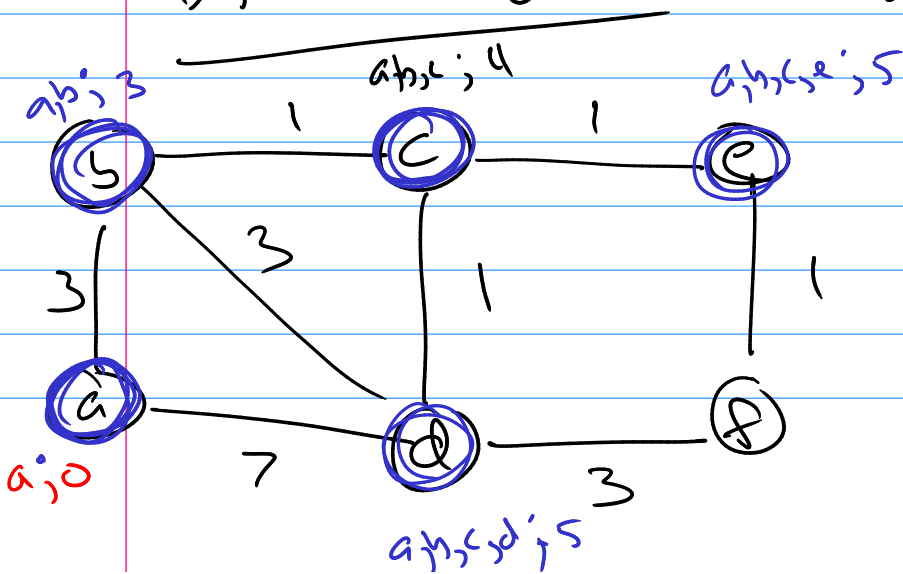
Thⁿ

The following statements are logically equivalent

- ① G is a tree (has no simple circuits)
- ② A graph is connected and if you remove any edge $(V, E - \{e\})$ that graph is disconnected.
- ③ A graph has no simple circuit, but if you add an edge it will.
- ④ A graph is connected and $|E| = |V| - 1$
- ⑤ Unique simple path between any two vertices.

Prax back Dijkstra's

Dijkstra's Algorithm



known
a,b

unknown
c,d,e,f

Vertex	Path	Cost
a	a	0
b	a,b	3
c	a,b,c	4
e	a,b,c,e	5
d	a,b,c,d	5
f	a,b,c,e,f	6

App

list of a to ()

list of b to ()

list of c to ()

;

last cost
~~paths~~

Middle?

Idea

for all least cost paths just count