

# Math 322

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Final Exam

4 probs / exam 1 to 3 so 12 probs

WSU Finals Schedule is

↳ Wed Dec 9<sup>th</sup>

1pm to 3pm

Our final is online so Wed Dec 9<sup>th</sup> 12pm - 4pm

Final is 12 problems taken from exam 1 to 3 / 2 hrs

Study? Solve Exam 1 to Exam 3!

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Q's

Key

0) Exam Start Time:

0) MyWSUId and Name:

1) Describe the language generated by the following expressions

a)  $0^*(101)1^*$

+5 all strings starting w/ any number of zero's contain 101 and end w/ any number of ones.

(ex) 0101

b)  $(0|1)^*111$

+5 Any string ending with 3 ones.

(ex) 0101010111

2) Give a regular expression that will generate the following languages

a) Any bit string with 111 in it.

+5  $(0|1)^* \underline{111} (0|1)^*$

b) Any bit strings that start with 01 and end with 101.

+5  $\underline{01} (0|1)^* \underline{101}$

Symbols | 01 is 001 is 011  
001 is 01

(ex) Start with 1 or more zeros

$0^n (0|1)^*$ ,  $n=1,2,3,\dots$

$00^* (0|1)^*$ ,  $0,1,2,3,\dots$

$$\tau = \epsilon$$

3) Name the grammar type (give its type number and name) and circle the productions that prevent it from being the next type.

a)  $S \rightarrow \underline{A}, S \rightarrow \underline{B}, S \rightarrow \lambda, A \rightarrow \underline{Sb}, B \rightarrow \underline{aB}, A \rightarrow a, \text{ and } B \rightarrow \underline{b}$

type 0, 1, 2, not 3 | Context free

b)  $S \rightarrow \underline{AB}, A \rightarrow aA, B \rightarrow bB, A \rightarrow a, \text{ and } B \rightarrow b$

type 2, not 3 | Context free

c)  $S \rightarrow AB, S \rightarrow \lambda, B \rightarrow aAb, \underline{aA \rightarrow a}, \text{ and } A \rightarrow B$

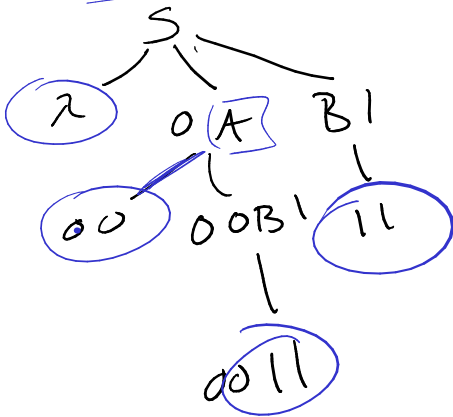
type 0, not 1 | Phase structure

d)  $\underline{S} \rightarrow \lambda, \underline{S} \rightarrow \underline{aA}, \underline{A} \rightarrow \underline{bB}, \underline{B} \rightarrow a, \text{ and } \underline{A} \rightarrow \underline{b}$

type 3 | regular grammar

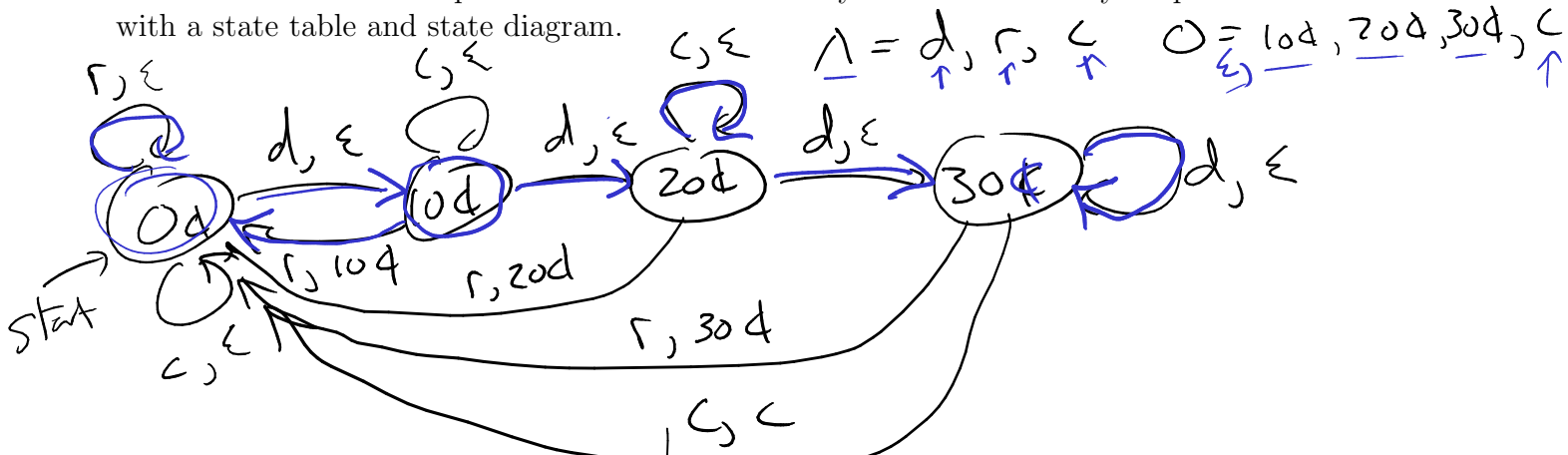
so phrase structure  
 not according to our book though.

4) For the grammar with  $V = \{0, 1, A, B, S\}$ ,  $T = \{0, 1\}$ , and the productions  $\underline{S} \rightarrow 0A$ ,  $S \rightarrow B1$ ,  $\underline{S} \rightarrow \lambda$ ,  $A \rightarrow 0B1$ ,  $B \rightarrow 1$ , and  $A \rightarrow 0$  find  $L(G)$ .



$$\underline{L(G)} = \lambda \mid 00 \mid 11 \mid 0011$$

5) Construct a finite-state machine with output that models a candy machine that accepts only dimes. Candy costs 30 cents and the machine keeps the money for any amount greater than 30 cents. The customer can push buttons to receive candy or to return money. Represent the machine with a state table and state diagram.



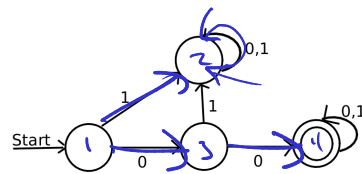
|     | outputs |     |   | transitions |    |     |
|-----|---------|-----|---|-------------|----|-----|
|     | d       | r   | c | d           | r  | c   |
| 0¢  | ε       | ε   | ε | 10¢         | 0¢ | 0¢  |
| 10¢ | ε       | 10¢ | ε | 20¢         | 0¢ | 10¢ |
| 20¢ | ε       | 20¢ | ε | 30¢         | 0¢ | 20¢ |
| 30¢ | ε       | 30¢ | ε | 30¢         | 0¢ | 0¢  |

6) For your machine in problem 5 what is the output and the state transitions if the customer's input string is "dime, return money, dime, dime, dime, dime, return money, dime, dime, candy"?

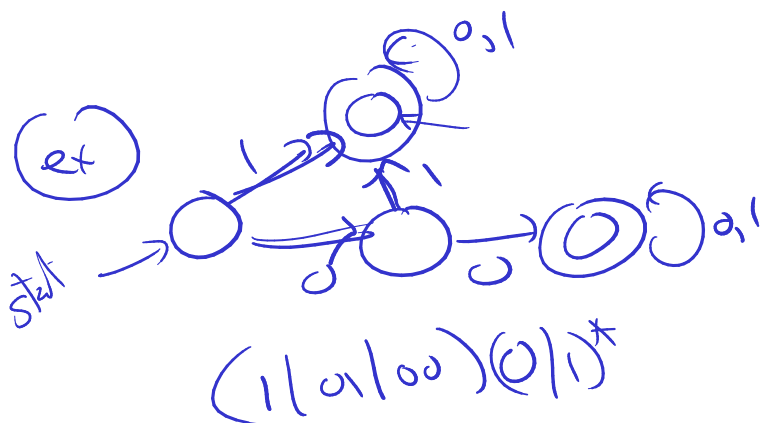
|        |     |     |     |     |     |     |     |     |     |     |
|--------|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|
| output | ε   | 10¢ | ε   | ε   | ε   | ε   | 30¢ | ε   | ε   | ε   |
| state  | 10¢ | 0¢  | 10¢ | 20¢ | 30¢ | 30¢ | 0¢  | 10¢ | 20¢ | 20¢ |

7) Determine the language recognized by a given deterministic finite-state automaton.

$00(011)^*$

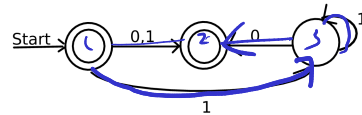


Final state:  $S_4$



8) Determine the language recognized by a given non-deterministic finite-state automaton.

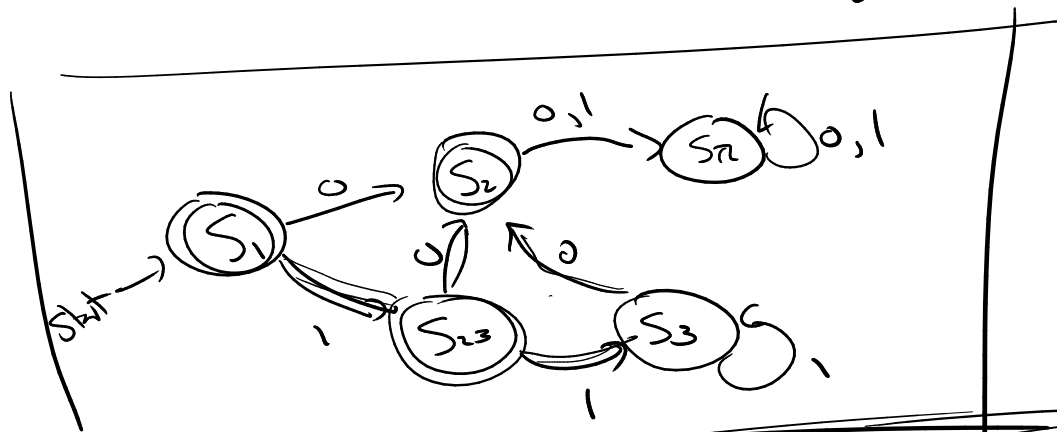
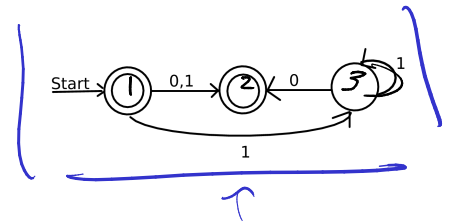
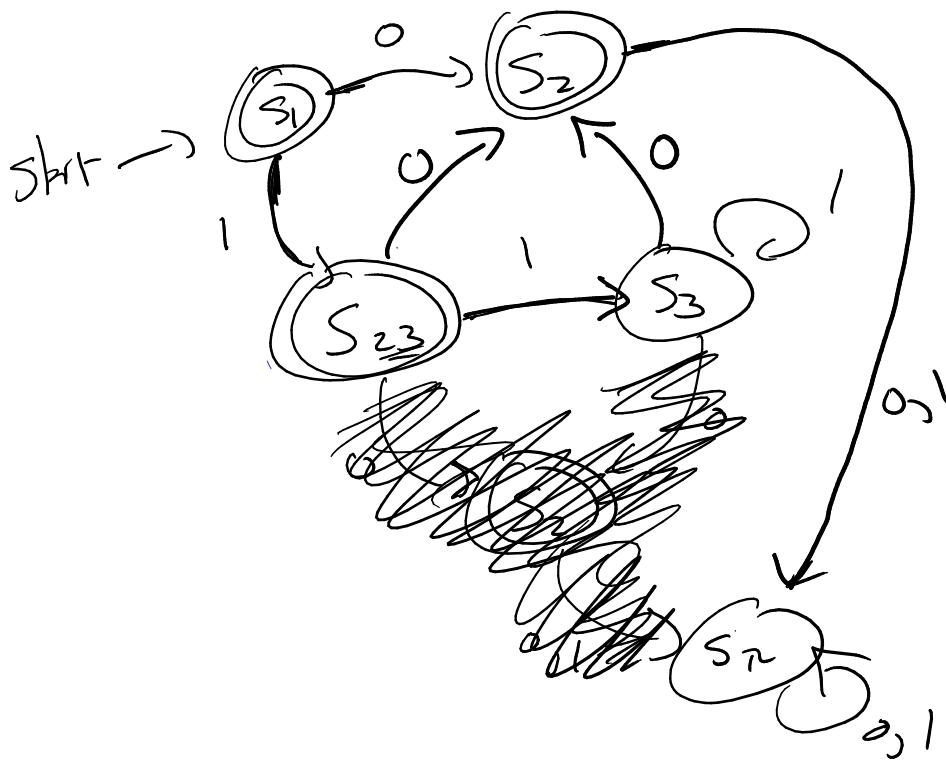
$S_1$   $\underline{\underline{\epsilon}}$  stay @ start  
 $S_2$  011 along top left



$11^*0$  from right

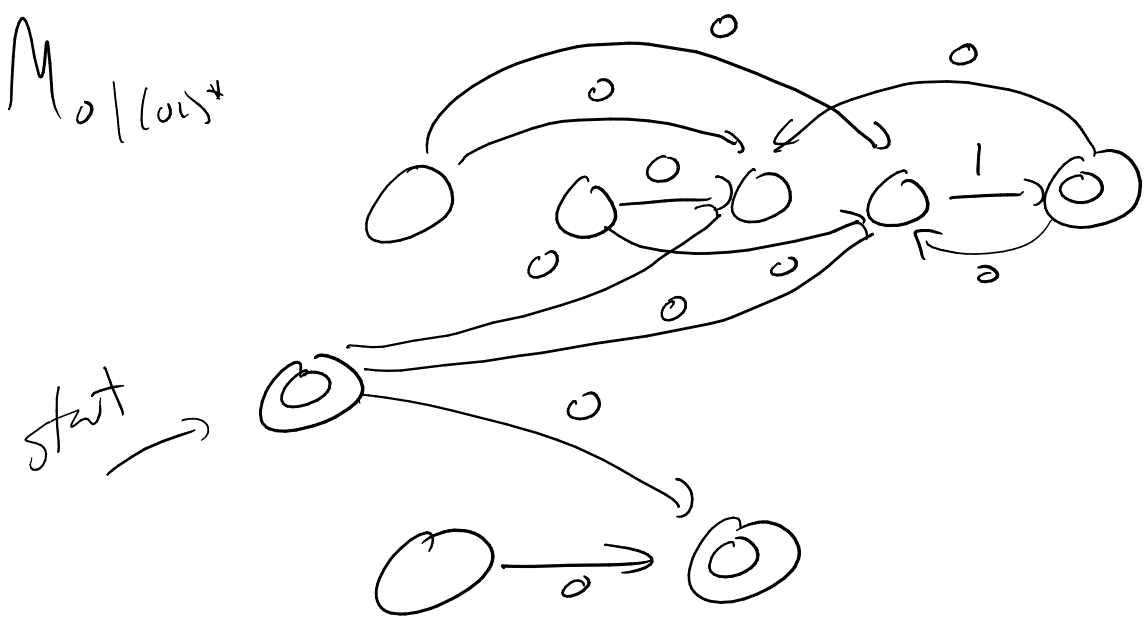
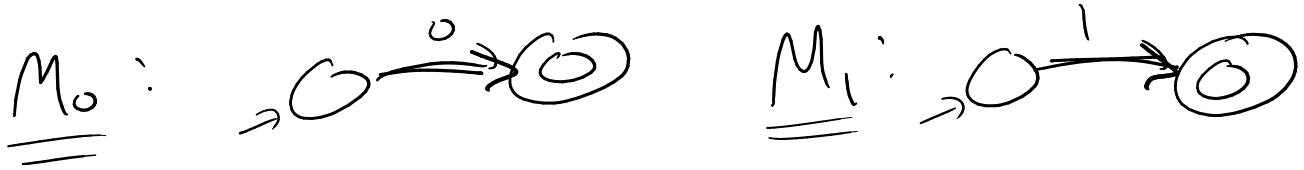
so  $\boxed{\epsilon \mid 0 \mid 1 \mid 11^*0} = L(M)$

9) For the non-deterministic finite-state automaton in problem 8 construct a deterministic finite-state automaton that recognizes the same language.



For NFA  
 $L(N) = \epsilon / 0 / 1 / 10 / 110 / 111^*0$   
check  $\epsilon / 0 / 1 / 111^*0$

10) Using the inductive constructions, find a non-deterministic finite-state automaton that recognizes  $0 \cup (01)^*$ .



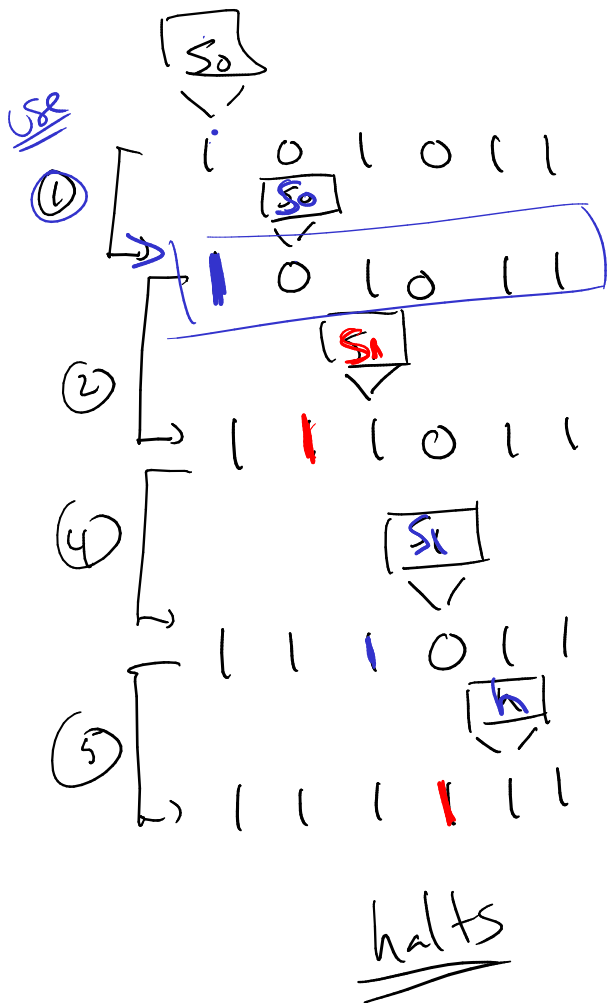
$$\textcircled{1} (s_0, 1) \rightarrow (1, R, s_0)$$

11) For the Turing machine defined by the five-tuples:

$(s_0, 1, 1, R, s_0)$ ,  $(s_0, 0, 1, R, s_1)$ ,  $(s_0, x, 1, R, s_1)$ ,  $(s_1, 1, 1, R, s_1)$ ,  $(s_1, 0, 1, R, h)$ , and  $(s_1, x, 1, R, h)$

Run the Turing machine on the below initial tape, write each of the positions, and determine the tape when T halts. What do you think T is doing?

Initial Tape: ... #, #, #, 1, 0, 1, 0, 1, 1, #, #, #, ...



$s_0$   $\equiv$  have not seen  $0$

$s_1$   $\equiv$  replaced 1<sup>st</sup>  $0$  w/  $1$

$h$   $\equiv$  replaced 2<sup>nd</sup>  $0$  w/  $1$   
or replaced blank after string w/  $1$ .



$n = 1, 2, 3, \dots$  ?

$n = 0, 1, 2, 3, \dots$  ?

$n \pmod 2 = 0$

12) Construct a Turing machine for  $f(n) = n \pmod 2$ . Run your machine on the input: 1,1,1,1,1.

1,1,1,1,1

0 is all blank tape  
 1 is 1  
 2 is 11  
 3 is 111  
 ...  
Positives

or  
 0 = 1  
 1 = 11  
 2 = 111  
 3 = 1111  
Non-neg's

if this is positive into it is 5. If non-neg into it is 4.

assume non-neg.

$S_0 \equiv$  have not "shared" yet (or back to start of sharing)

$S_1 \equiv$  one given out  $S_0$

$(S_0, 1) \rightarrow (\#, R, S_1)$

$(S_1, 1) \rightarrow (\#, R, S_0)$

$(S_0, \#) \rightarrow (1, R, S_1)$

$(S_1, \#) \rightarrow (1, R, S_0)$

|   |   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|---|
| 4 =   | 1 | 1 | 1 | 1 | 1 | # | # | # |
| 3 =   | # | 1 | 1 | 1 | 1 | # | # | # |
| 2 =   | # | # | 1 | 1 | 1 | # | # | # |
| 1 =   | # | # | # | 1 | 1 | # | # | # |
| 0 =   | # | # | # | # | 1 | # | # | # |
| <span style="border: 1px solid black; padding: 2px;"> </span> = | # | # | # | # | # | # | # | # |
| 0 =   | # | # | # | # | # | 1 | # | # |

$n \pmod 2 = 0$