

Math 530

Q's

2.1 2

Euler circuits
(need all even degree)

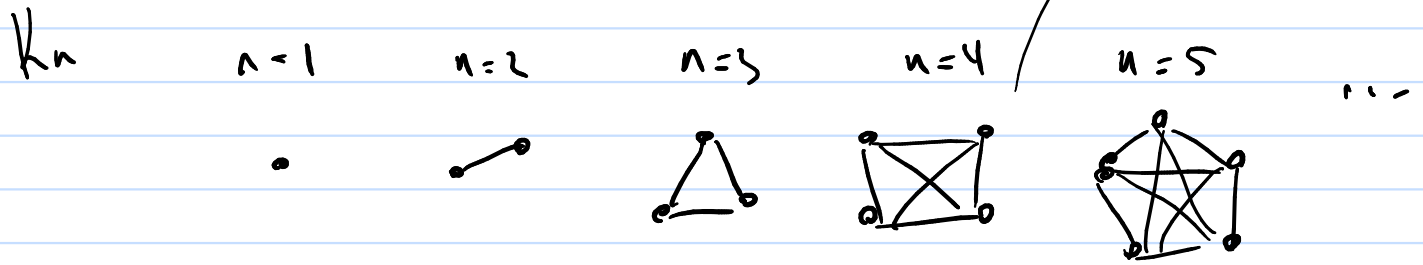
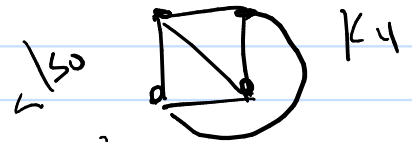
Note

1.4 (1, 3, b, d, f, 7, a, c, d)

2.1 (1, 2, 3, 9)

2.2 (1, 3, 5, 7, a)

Euler paths
(exactly 2 of odd degree)



$ V $	1	2	3	4	5	n
$ E $	0	1	3	6	10	$\frac{n(n-1)}{2}$
$\deg(v)$	0	1	2	3	4	$(n-1)$

2a) \cup $\deg(v)$ is all evens so $|V| = 3, 5, 7, \dots$

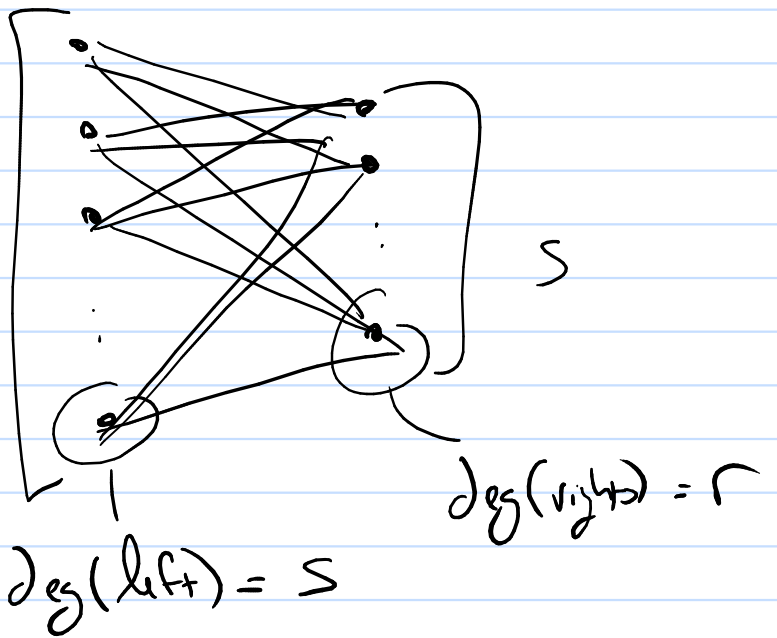
2b) K_2 it is a, b

2c) $K_{r,s}$

$2c$ $K_{r,s}$

Euler cycles = all even

So r, s are even.



Euler path \Rightarrow

if one side is odd

the other has only 2 vertices

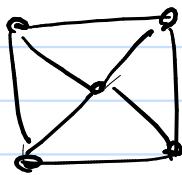
$K_{2, \text{odd}}$ has Euler paths

(*) W_n

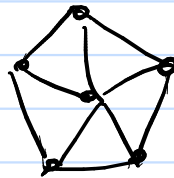
$n=3$



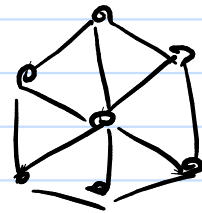
$n=4$



$n=5$



$n=6$



n

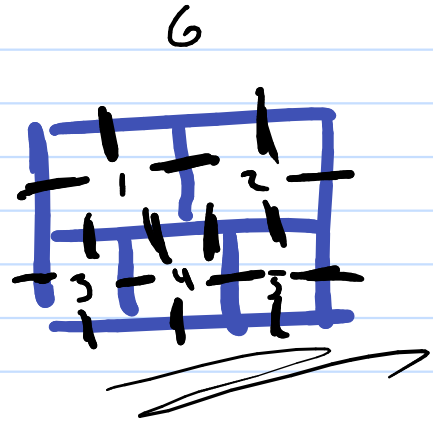
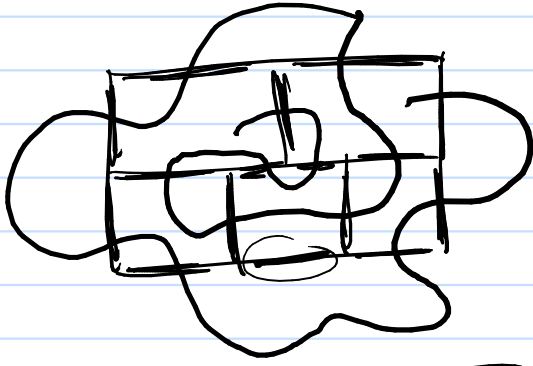
	$n=3$	$n=4$	$n=5$	$n=6$	n
$ V $	4	5	6	7	$(n+1)$
$ E $	6	8	10	12	$2n$
\deg_{outer}	3	4	5	6	n
\deg_{center}	3	3	3	3	3

No Euler path

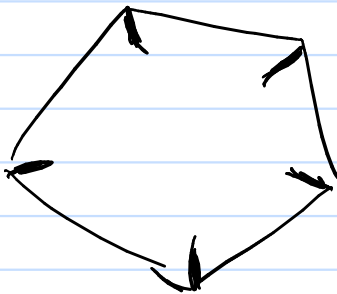
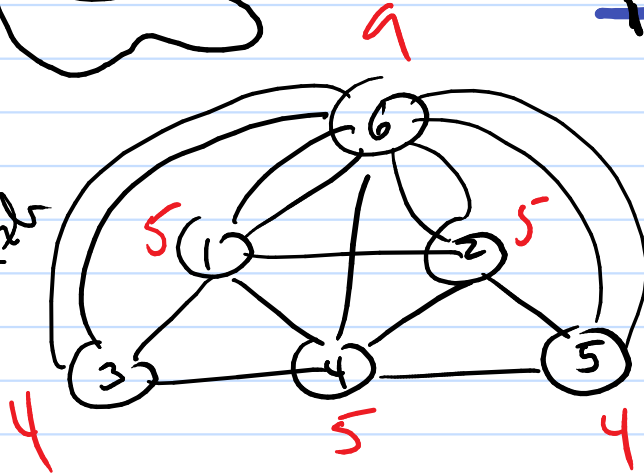
no Euler circuit

Apps of Euler Problems

①
cut
puzzle



↖ dual
cut puzzle

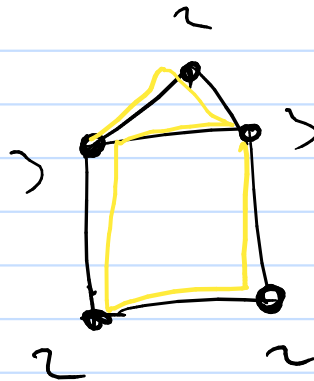


deg of outside
= 5



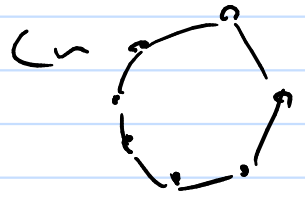
tracing

draw and
do not stop



" Euler path or Euler circuit

Hamiltonian:



Dirac

if $\deg(v) \geq \frac{|V|}{2}$ then
 G has a Hamiltonian Circuit

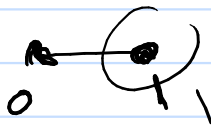
Ore's

if $\deg(v_i) + \deg(v_j) \geq |V|$
with v_i, v_j are non-adjacent then
 G has a Hamiltonian Circuit

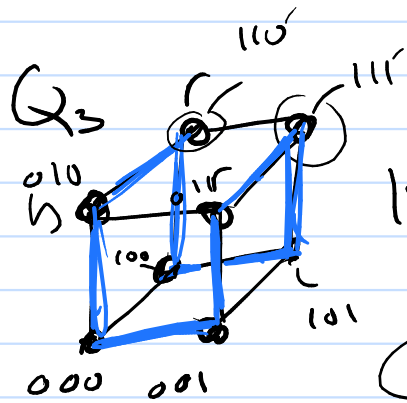
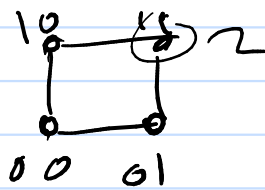
Q_0



Q_1



Q_2



$|V| = 8$

000 = 0₁

001

010

011

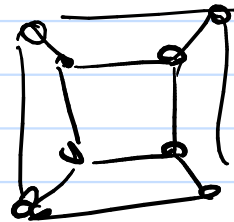
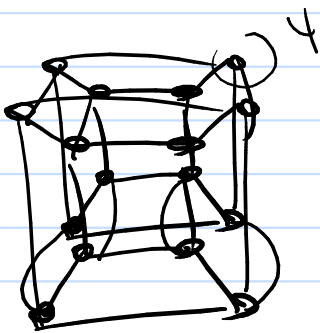
100

101

110

111 = 0₂

Q_4

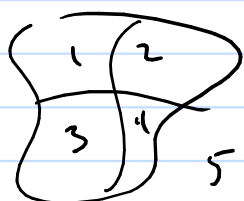


2.3

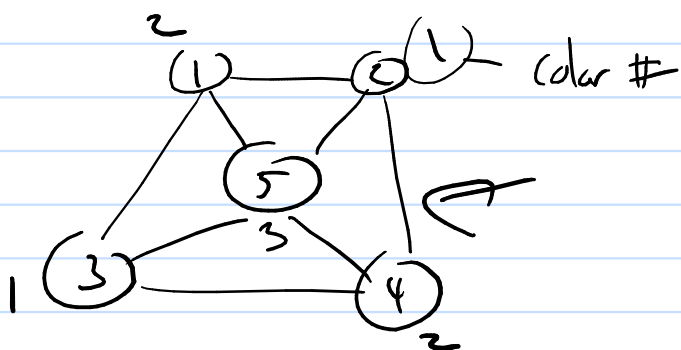
Graph coloring

adj. vertices have different colors

(ex)



dual



chromatic number of a graph is the absolute min. number of colors needed to color the graph

Note

to show the chromatic number is K means to

show \nexists $K-1$ coloring for a graph

has two adj. vertices w/ same color.

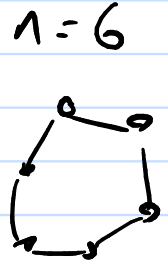
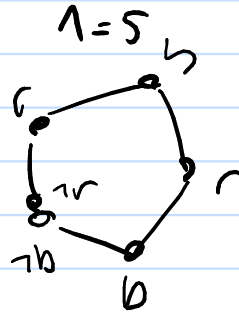
this is a NP-complete problem

Facts:

- (1) if G is planar, chromatic number is ≤ 4
- (2) K_n has a chromatic number of n .
- (3) bipartite has a chromatic number of 2.
- (4) if G has a K_i subgraph for i of its vertices then i is a lower bound for its chromatic number.

(ex) C_n

$n=3$



chromatic number

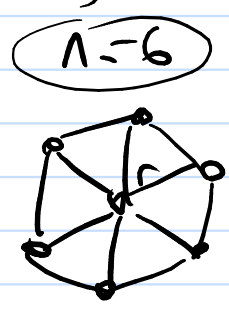
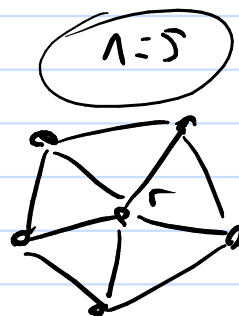
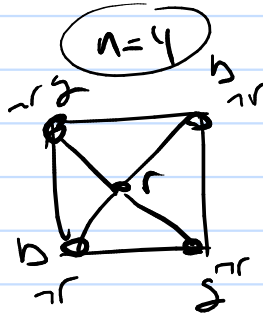
3

2

3

2

W_n



chromatic

4

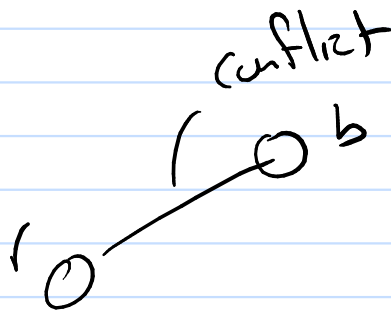
3

4

3

Applicaz:

- Scheduling



chromatic number = number of isolated locations